

Concurrent programming: From theory to practice

Concurrent Algorithms 2018

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From theory to practice

Theoretical
(design)

Practical
(design)

Practical
(implementation)

From theory to practice

Theoretical
(design)

Practical
(design)

Practical
(implementation)

- Impossibilities
- Upper/Lower bounds
- Techniques
- System models
- Correctness proofs



**Design
(pseudo-code)**

From theory to practice

Theoretical (design)

- Impossibilities
- Upper/Lower bounds
- Techniques
- System models
- Correctness proofs



**Design
(pseudo-code)**

Practical (design)

- System models
 - shared memory
 - message passing
- **Finite memory**
- Practicality issues
 - re-usable objects
- **Performance**



**Design
(pseudo-code,
prototype)**

Practical (implementation)

From theory to practice

Theoretical (design)

- Impossibilities
- Upper/Lower bounds
- Techniques
- System models
- Correctness proofs



**Design
(pseudo-code)**

Practical (design)

- System models
 - shared memory
 - message passing
- **Finite memory**
- Practicality issues
 - re-usable objects
- **Performance**



**Design
(pseudo-code,
prototype)**

Practical (implementation)

- **Hardware**
- Which atomic ops
- Memory consistency
- Cache coherence
- Locality
- **Performance**
- **Scalability**



**Implementation
(code)**

Outline

- CPU caches
- Cache coherence
- Placement of data
- Graph processing: Concurrent data structures

Outline

- **CPU caches**
- Cache coherence
- Placement of data
- Graph processing: Concurrent data structures

Why do we use caching?

Core

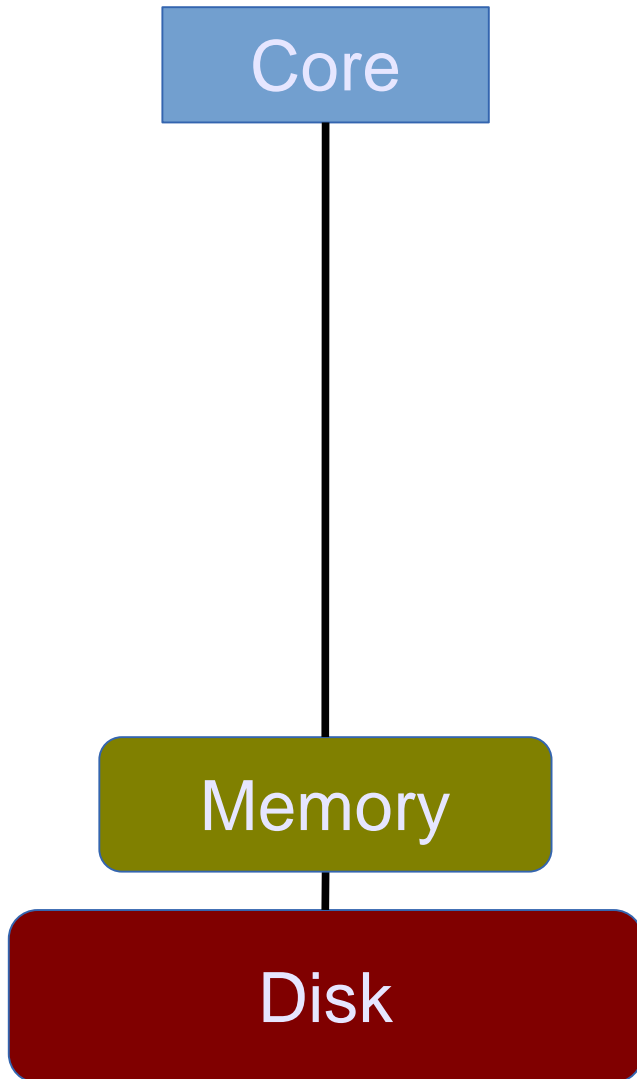


- Core freq: 2GHz = 0.5 ns / instr
- Core → Disk = ~ms

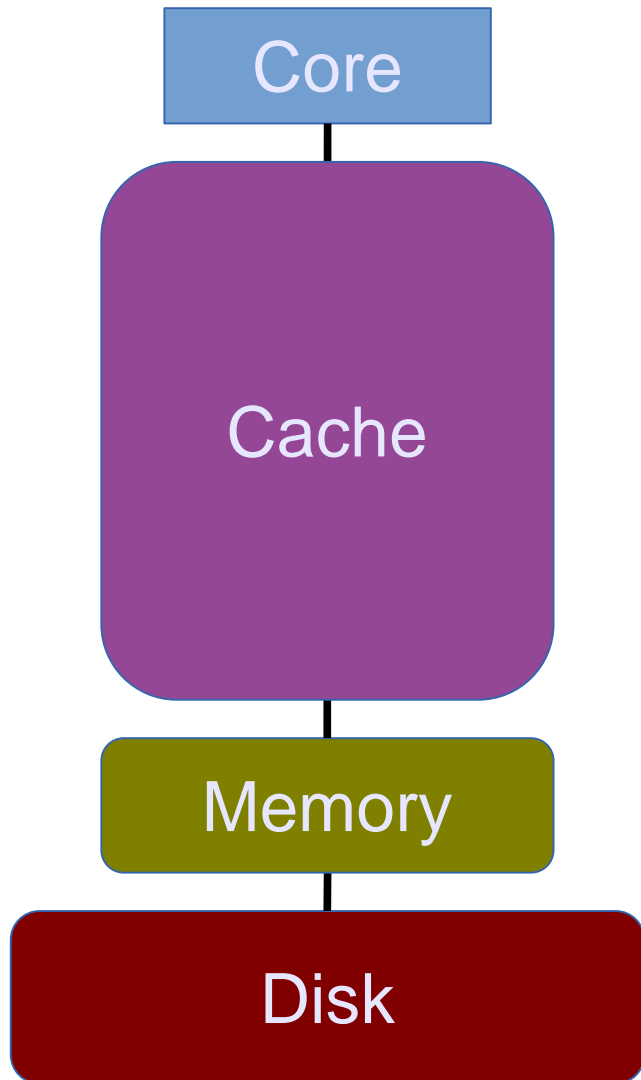
Disk

Why do we use caching?

- Core freq: 2GHz = 0.5 ns / instr
- Core → Disk = ~ms
- Core → Memory = ~100ns

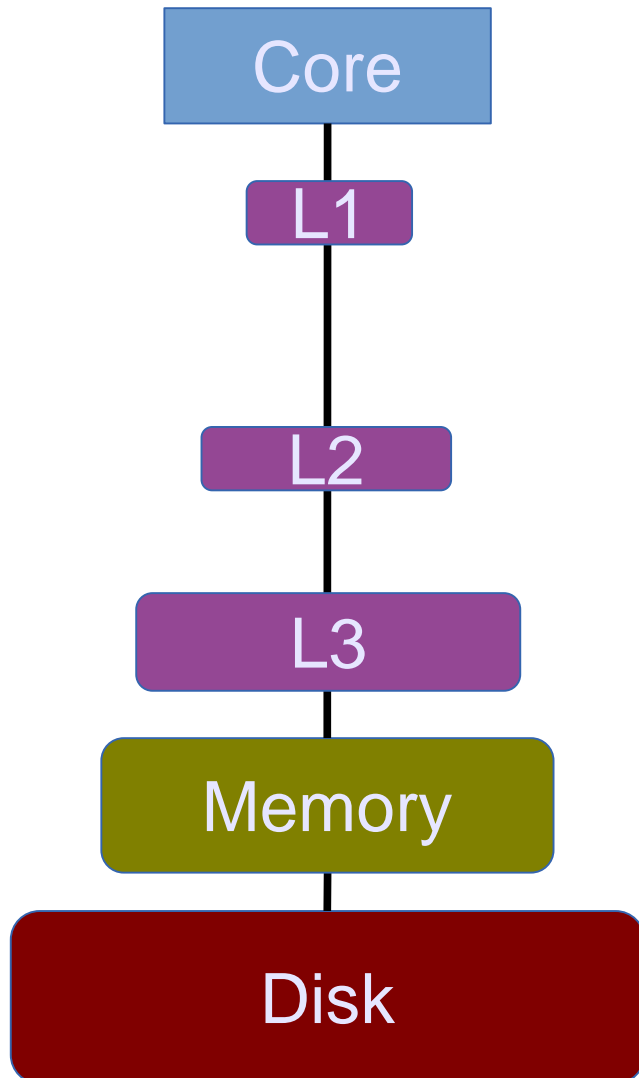


Why do we use caching?



- Core freq: 2GHz = 0.5 ns / instr
- Core → Disk = ~ms
- Core → Memory = ~100ns
- **Cache**
 - Large = slow
 - Medium = medium
 - Small = fast

Why do we use caching?

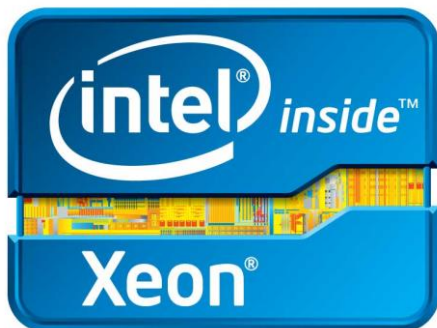


- Core freq: 2GHz = 0.5 ns / instr
- Core → Disk = ~ms
- Core → Memory = ~100ns
- Cache
 - Core → L3 = ~20ns
 - Core → L2 = ~7ns
 - Core → L1 = ~1ns

Typical server configurations

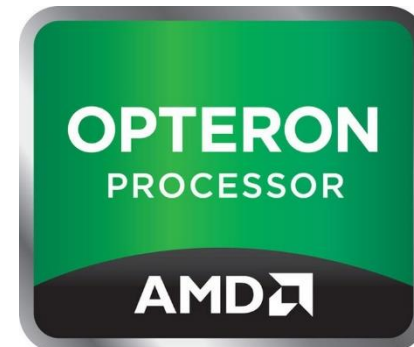
- **Intel Xeon**

- 12 cores @ 2.4GHz
- L1: 32KB
- L2: 256KB
- L3: 40MB
- Memory: 128GB



- **AMD Opteron**

- 12 cores @ 2.4GHz
- L1: 64KB
- L2: 512KB
- L3: 20MB
- Memory: 128GB



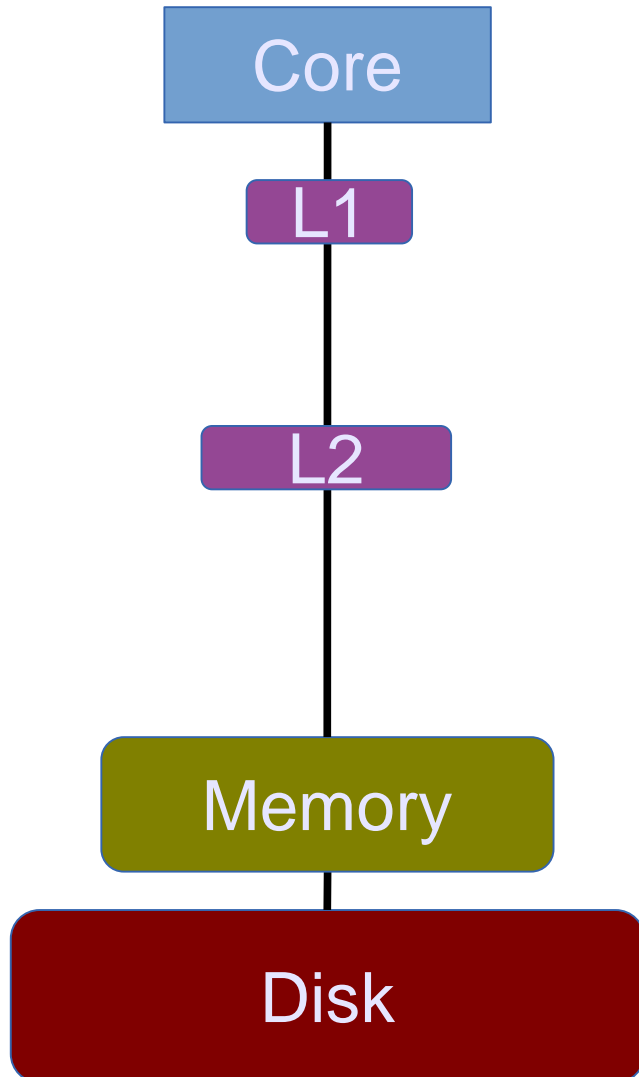
Experiment

Throughput of accessing some memory,
depending on the memory size

Outline

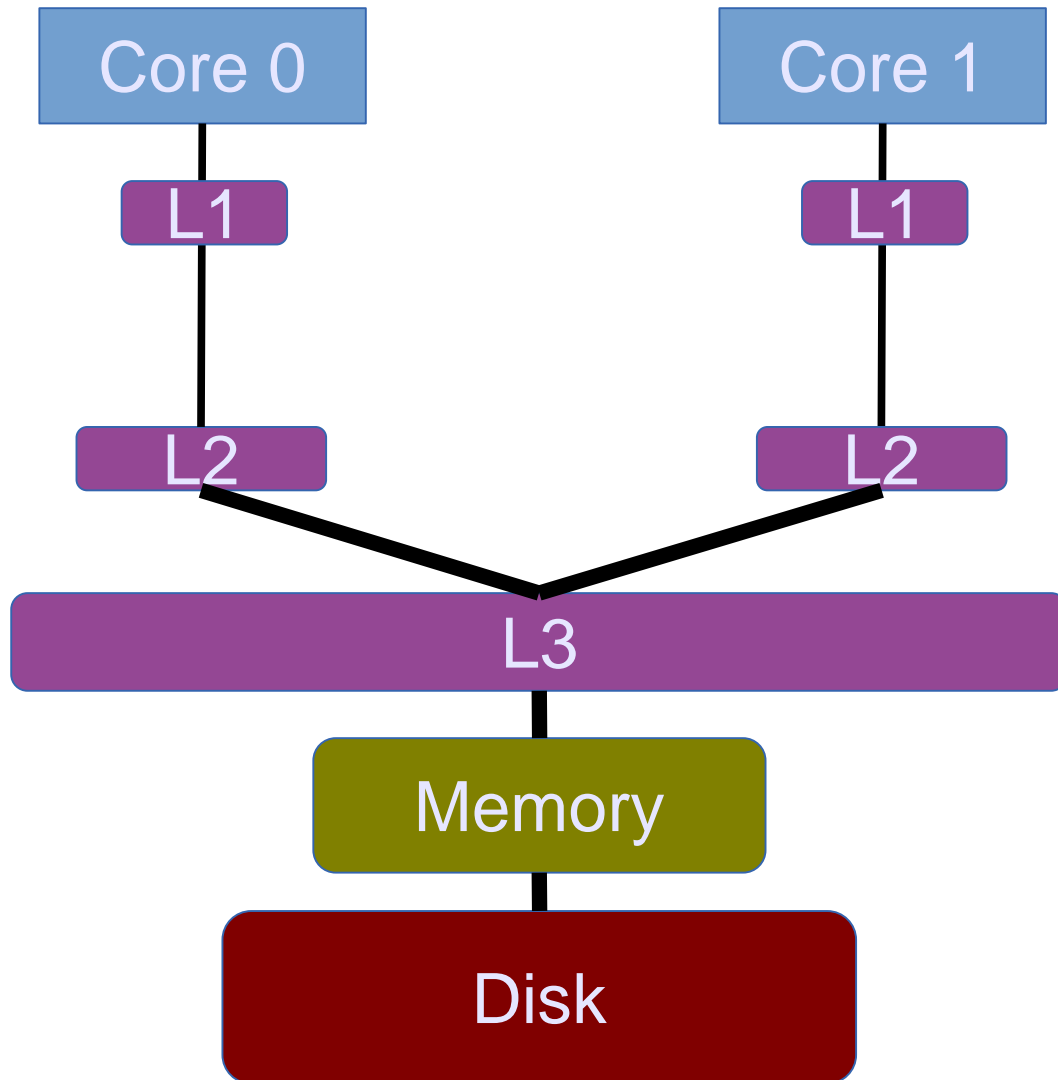
- CPU caches
- **Cache coherence**
- Placement of data
- Graph processing: Concurrent data structures

Until ~2004: single-cores



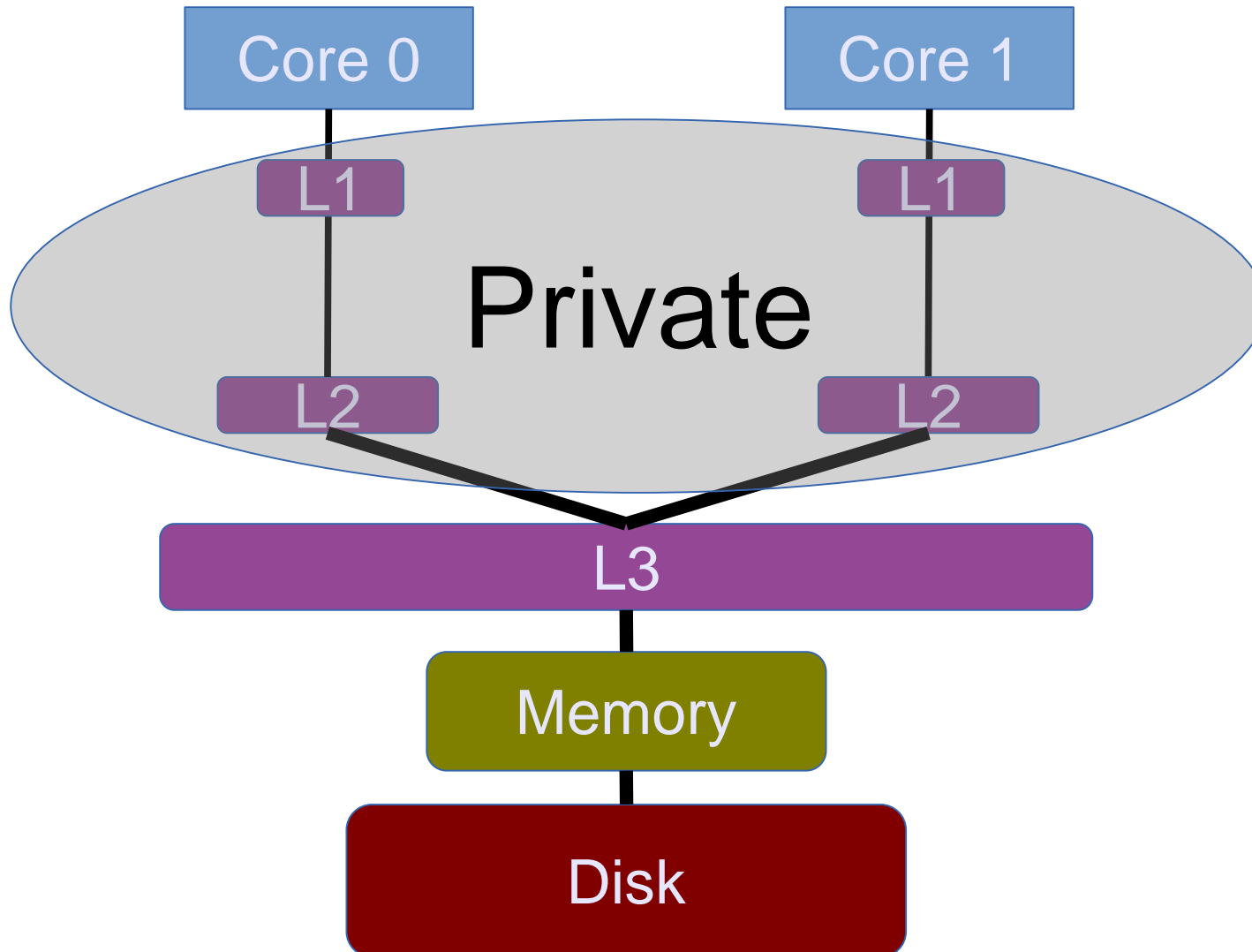
- Core freq: 3+GHz
- Core → Disk
- Core → Memory
- Cache
 - Core → L3
 - Core → L2
 - Core → L1

After ~2004: multi-cores



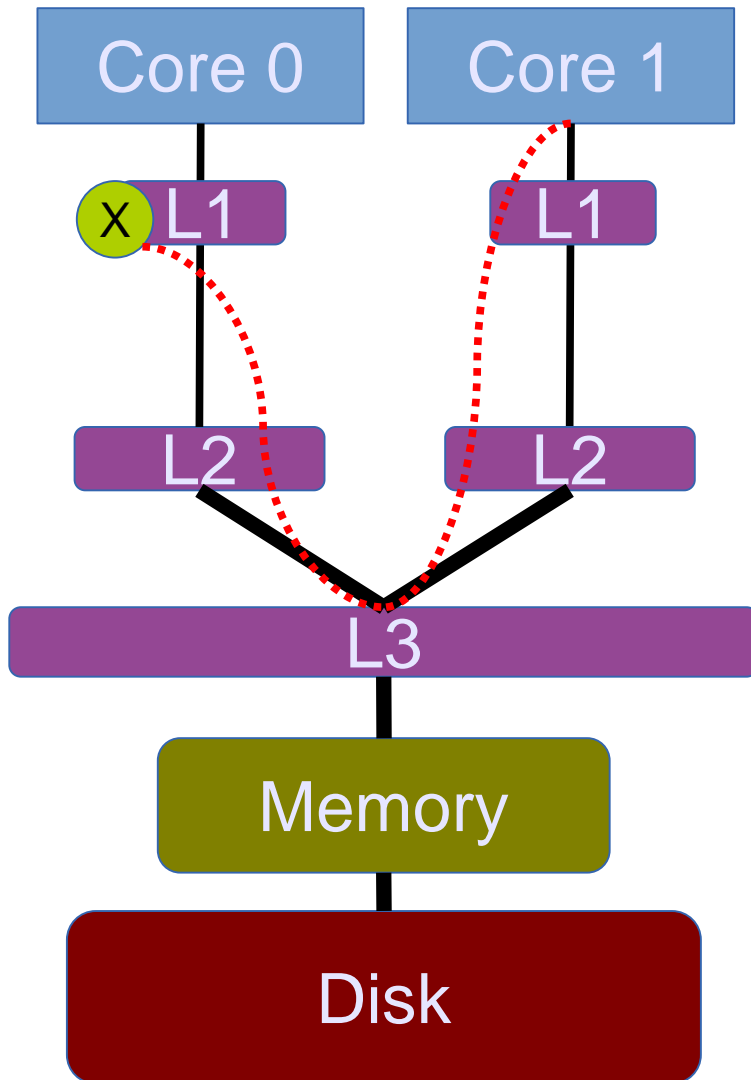
- Core freq: ~2GHz
- Core → Disk
- Core → Memory
- Cache
 - Core → **shared** L3
 - Core → L2
 - Core → L1

Multi-cores with private caches



=
multiple
copies

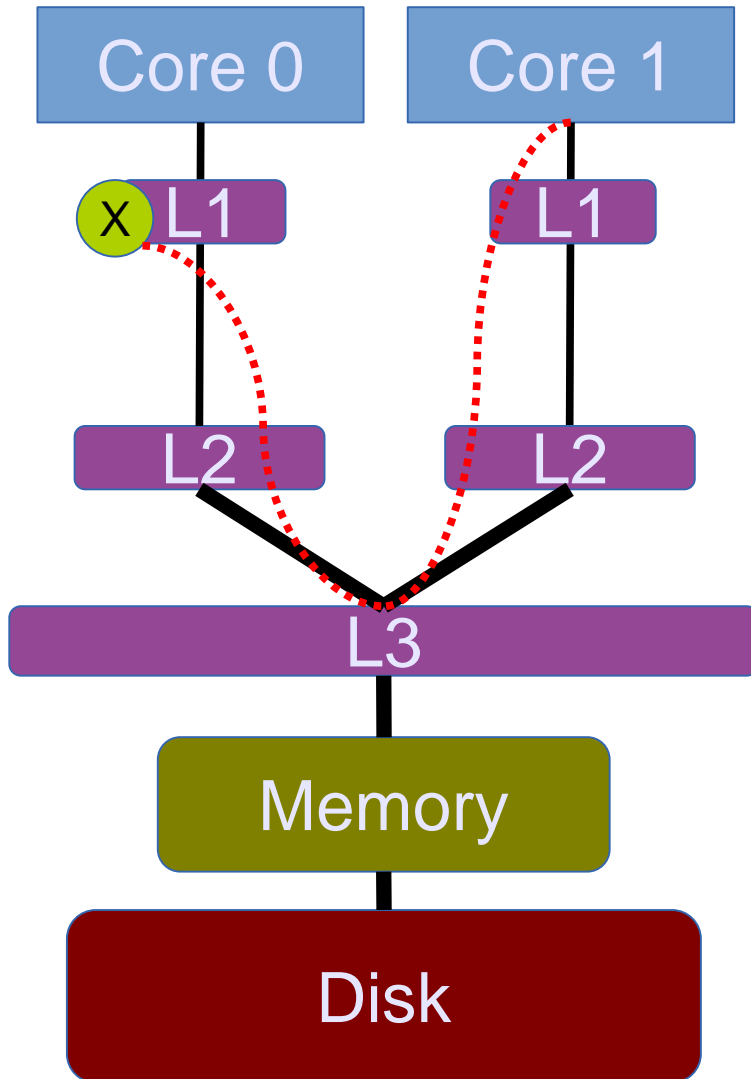
Cache coherence for consistency



Core 0 has **X** and Core 1

- wants to write on **X**
- wants to read **X**
- did Core 0 write or read **X**?

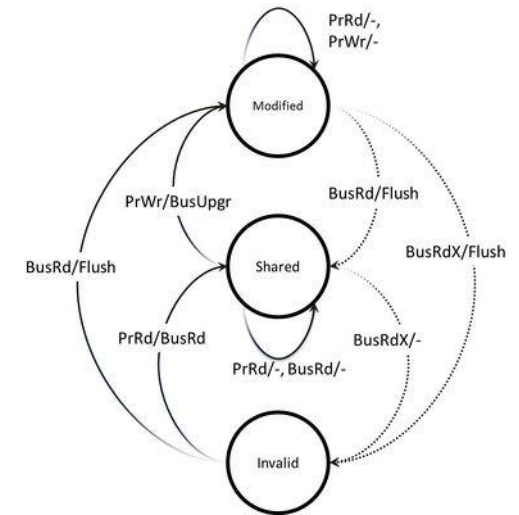
Cache coherence principles



- To perform a **write**
 - invalidate all readers, or
 - previous writer
- To perform a **read**
 - find the latest copy

Cache coherence with MESI

- A state diagram
- State (per cache line)
 - **Modified**: the only dirty copy
 - **Exclusive**: the only clean copy
 - **Shared**: a clean copy
 - **Invalid**: useless data



The ultimate goal for scalability

- Possible states
 - **Modified**: the only dirty copy
 - **Exclusive**: the only clean copy
 - **Shared**: a clean copy
 - **Invalid**: useless data
- **Which state is our “favorite”?**

The ultimate goal for scalability

- Possible states

- **Modified**: the only dirty copy
- **Exclusive**: the only clean copy

- **Shared**: a clean copy

- **Invalid**: useless data

= threads can keep the data close (L1 cache)

= faster

Experiment

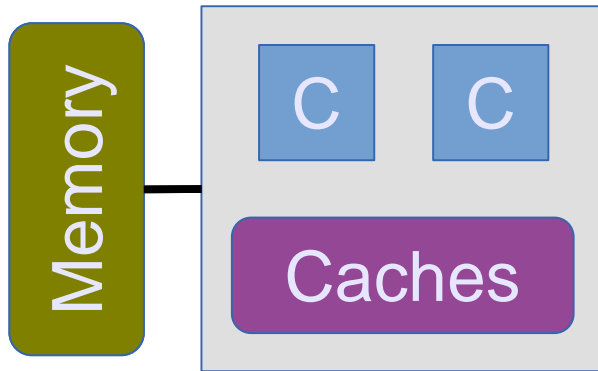
The effects of false sharing

Outline

- CPU caches
- Cache coherence
- **Placement of data**
- Graph processing: Concurrent data structures

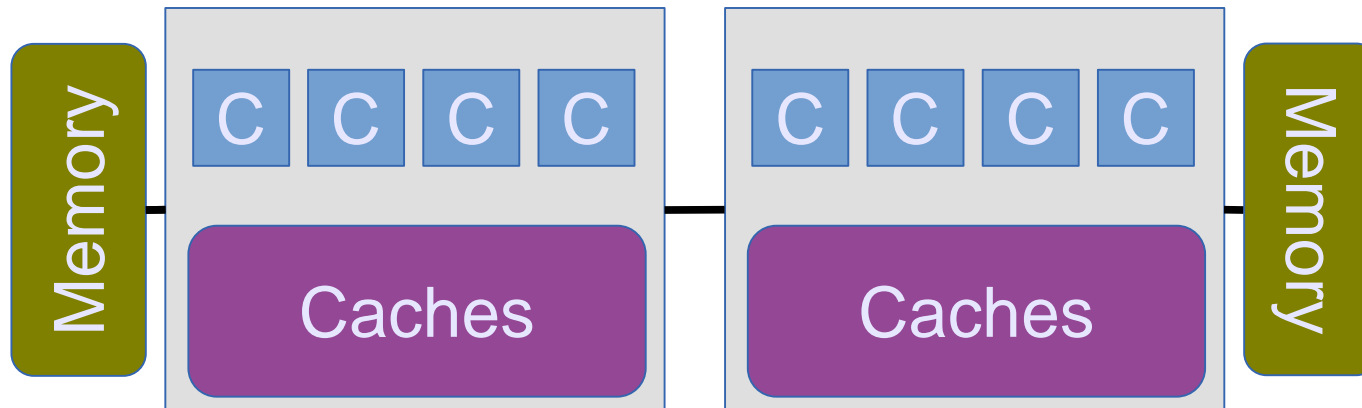
Uniformity vs. non-uniformity

- Typical **desktop** machine



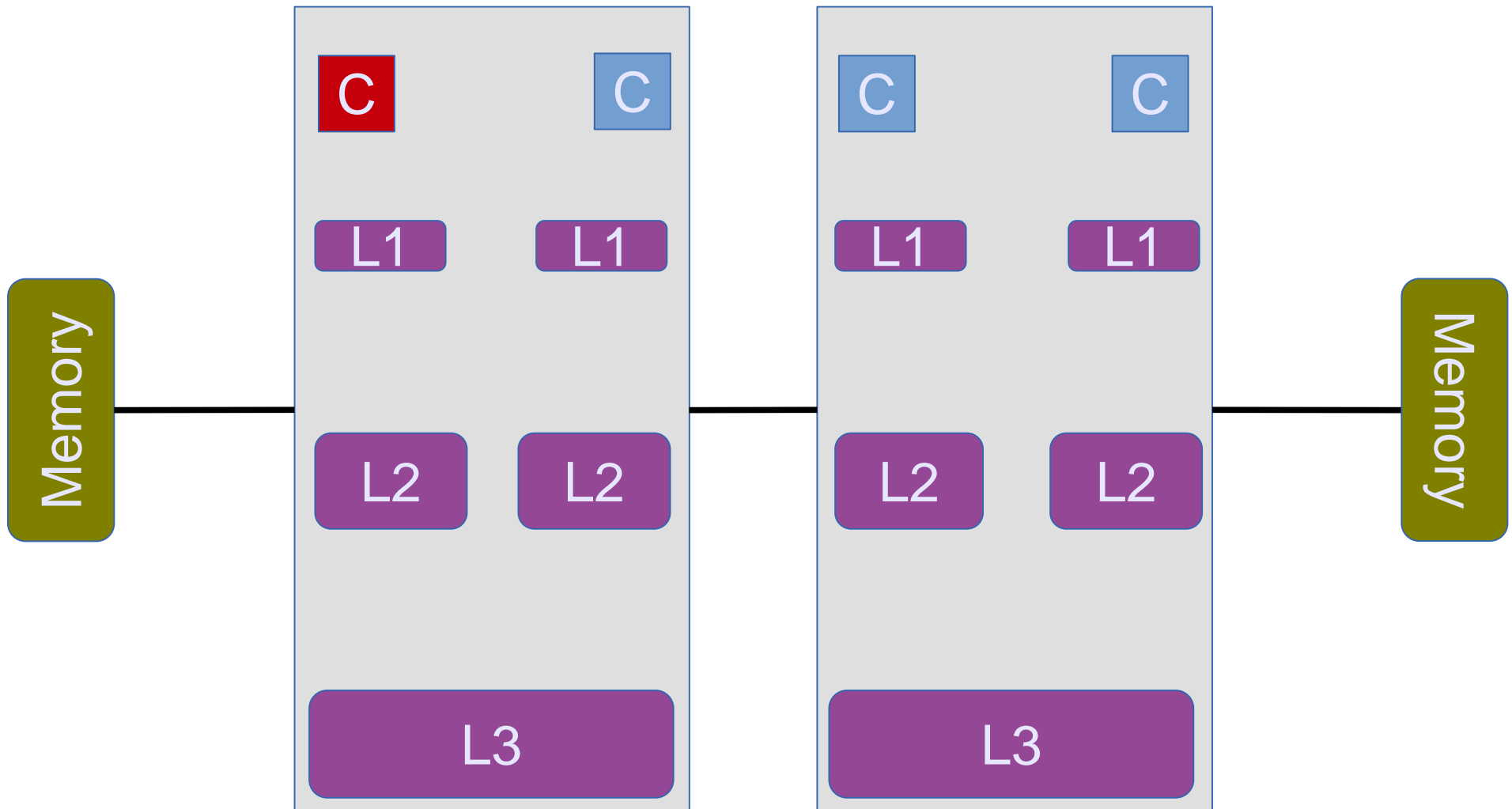
= Uniform

- Typical **server** machine

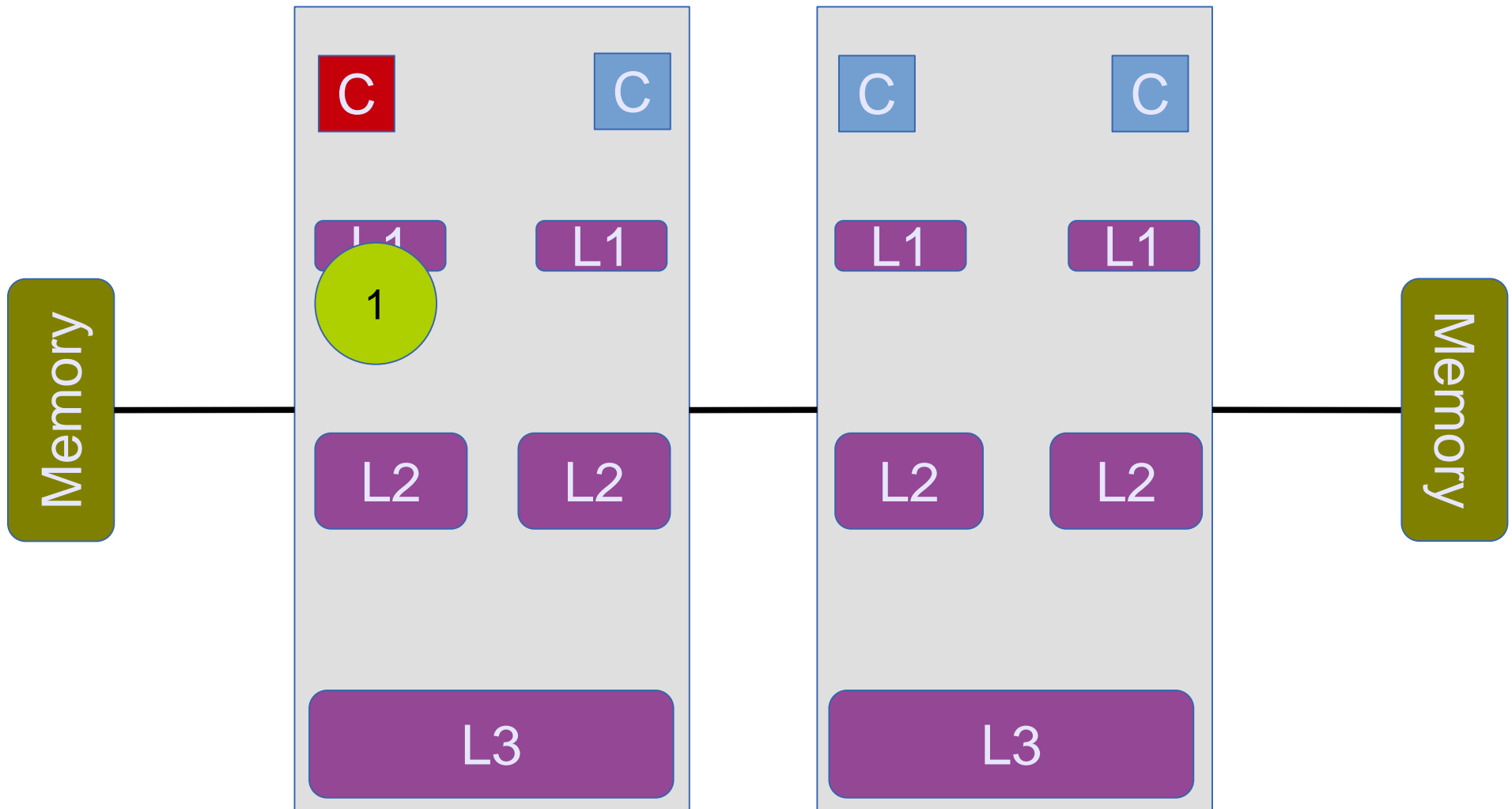


= non-Uniform

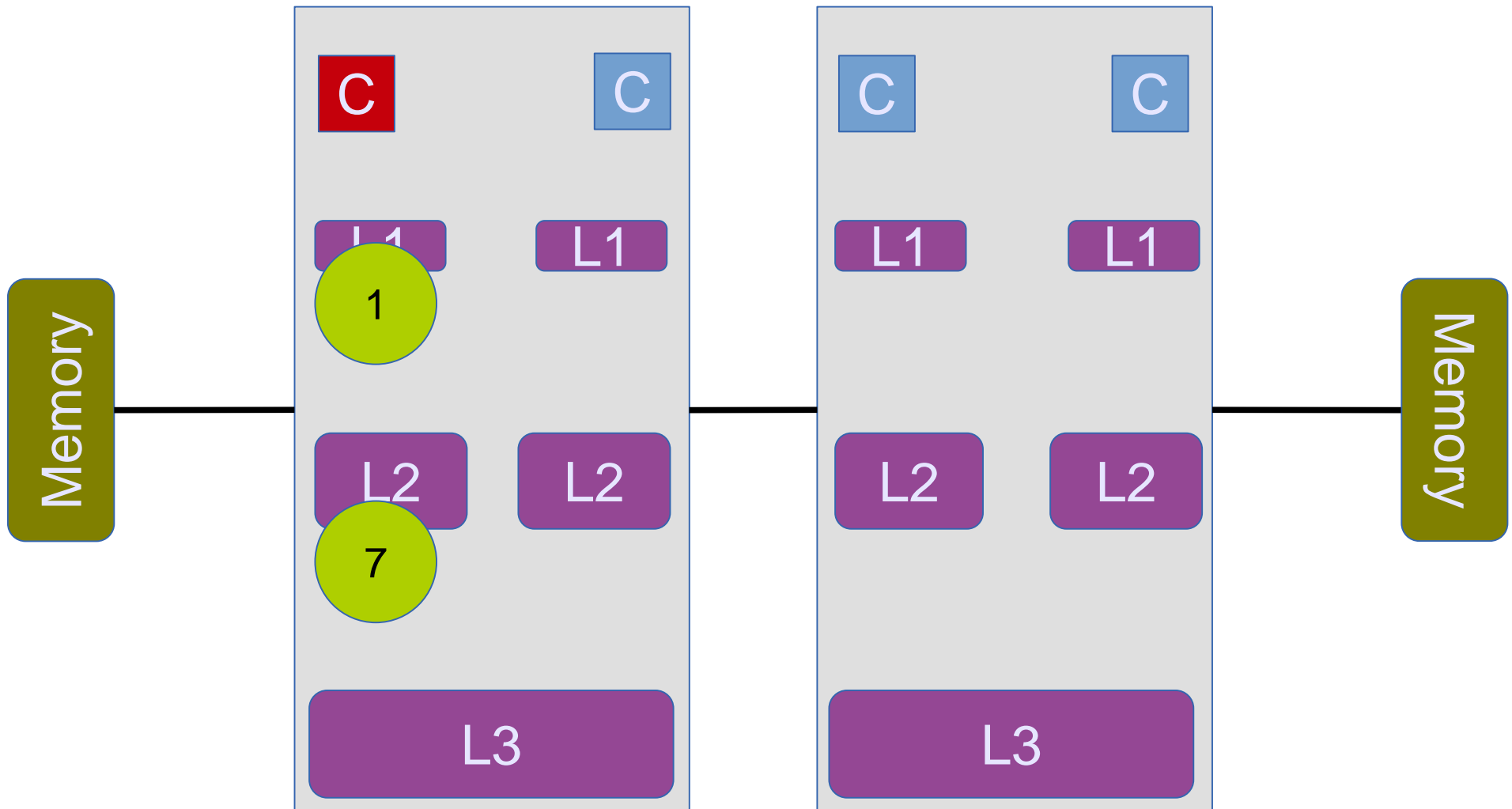
Latency (ns) to access data



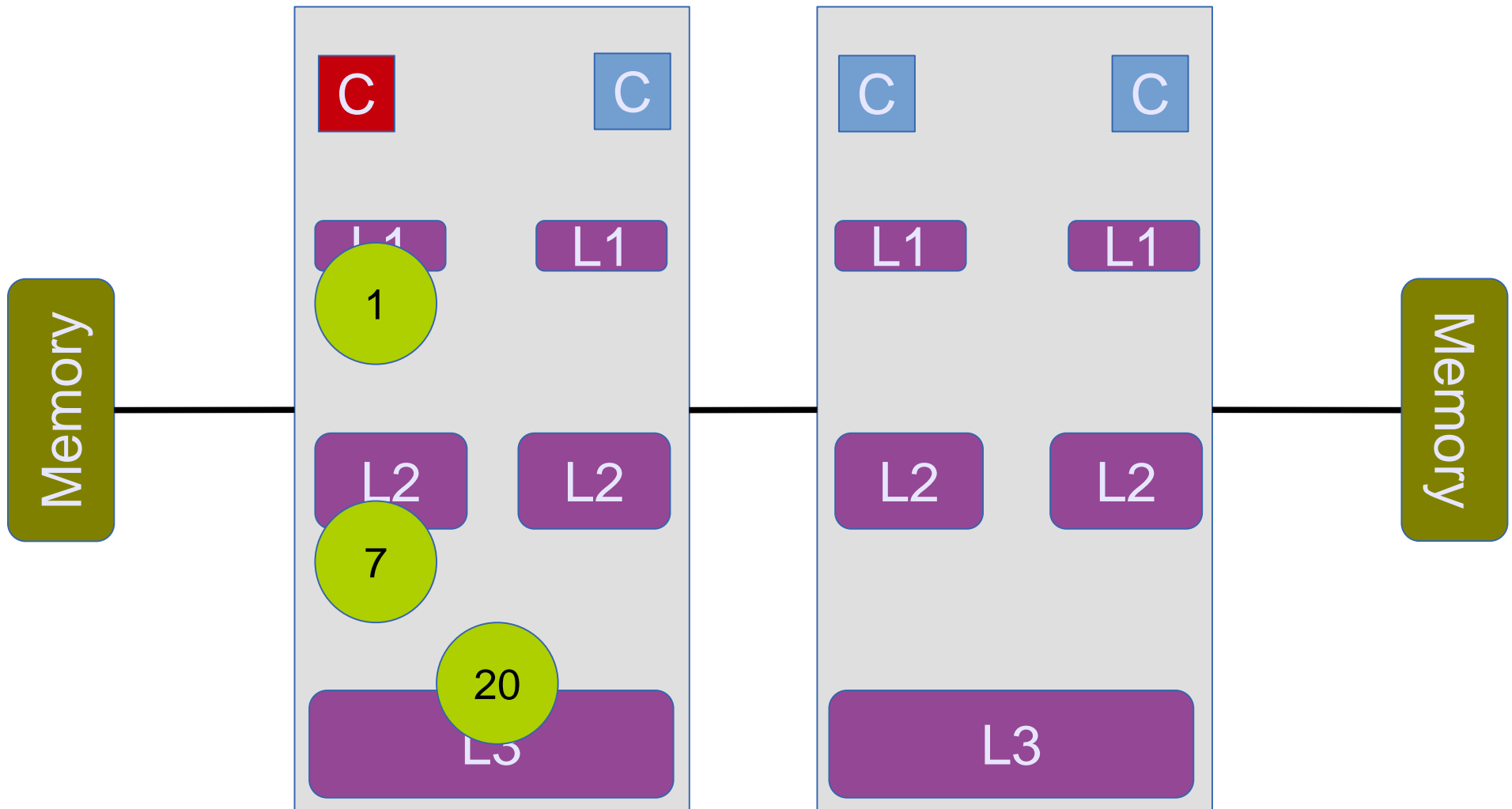
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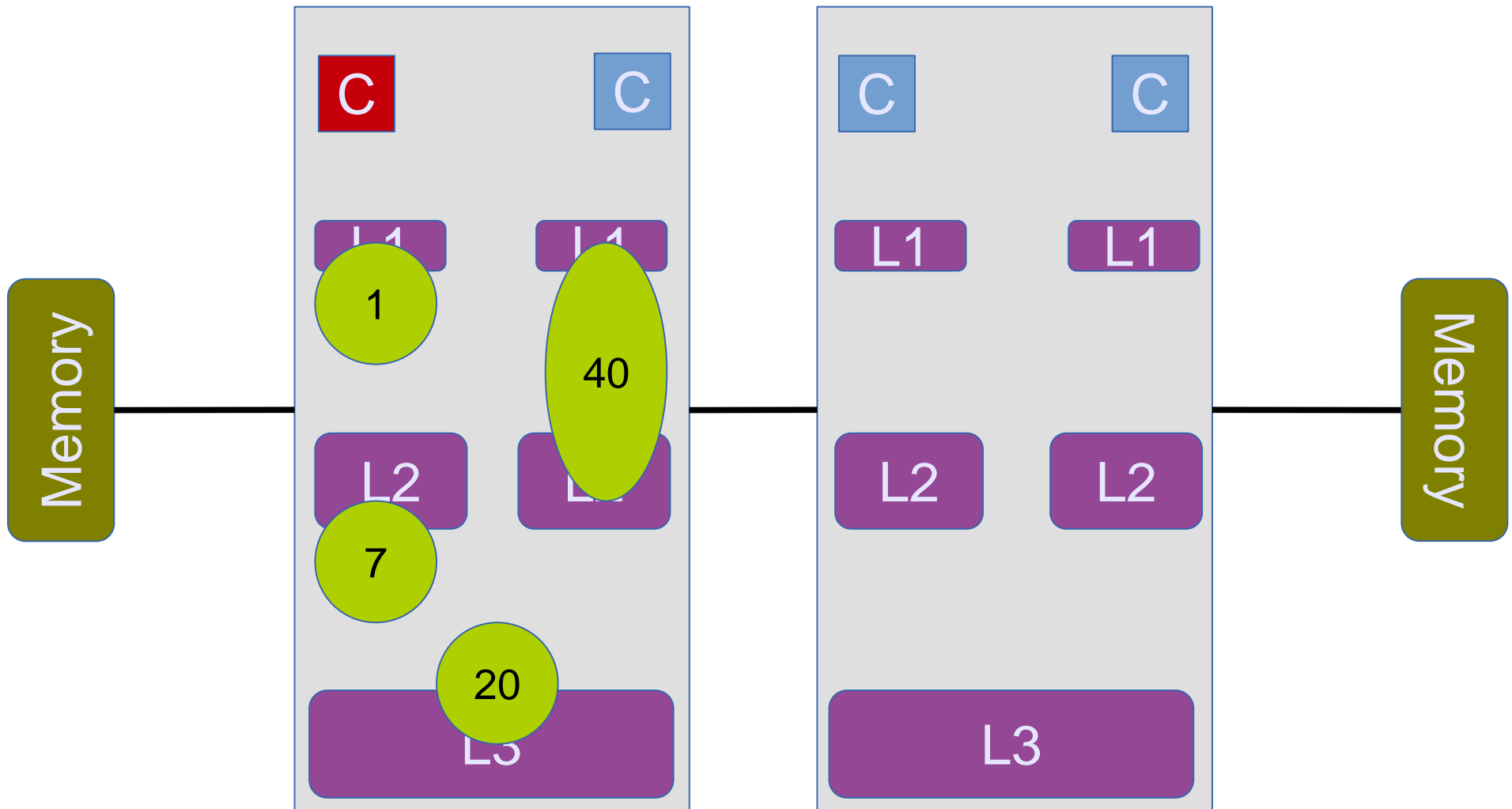
Latency (ns) to access data



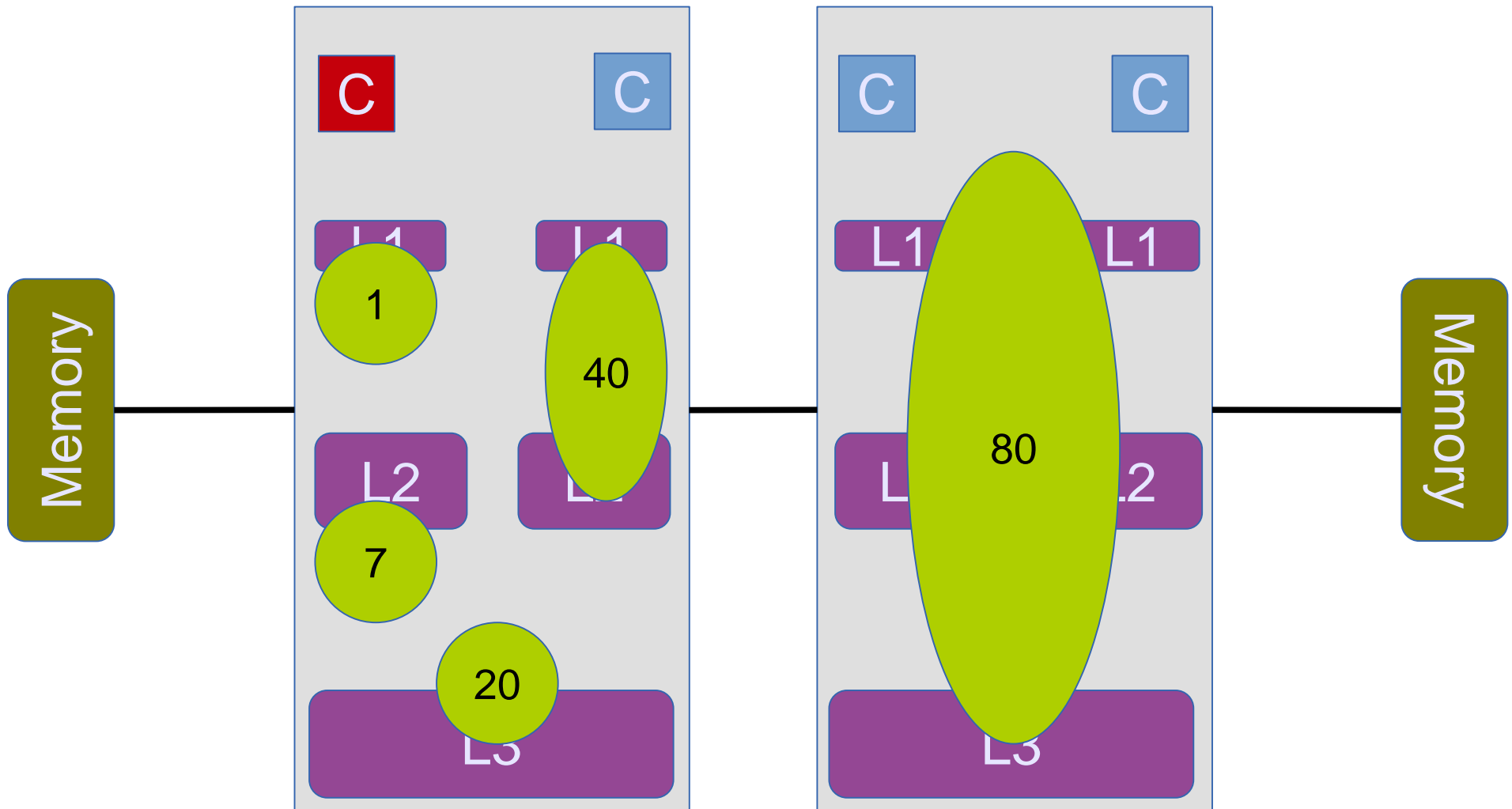
Latency (ns) to access data



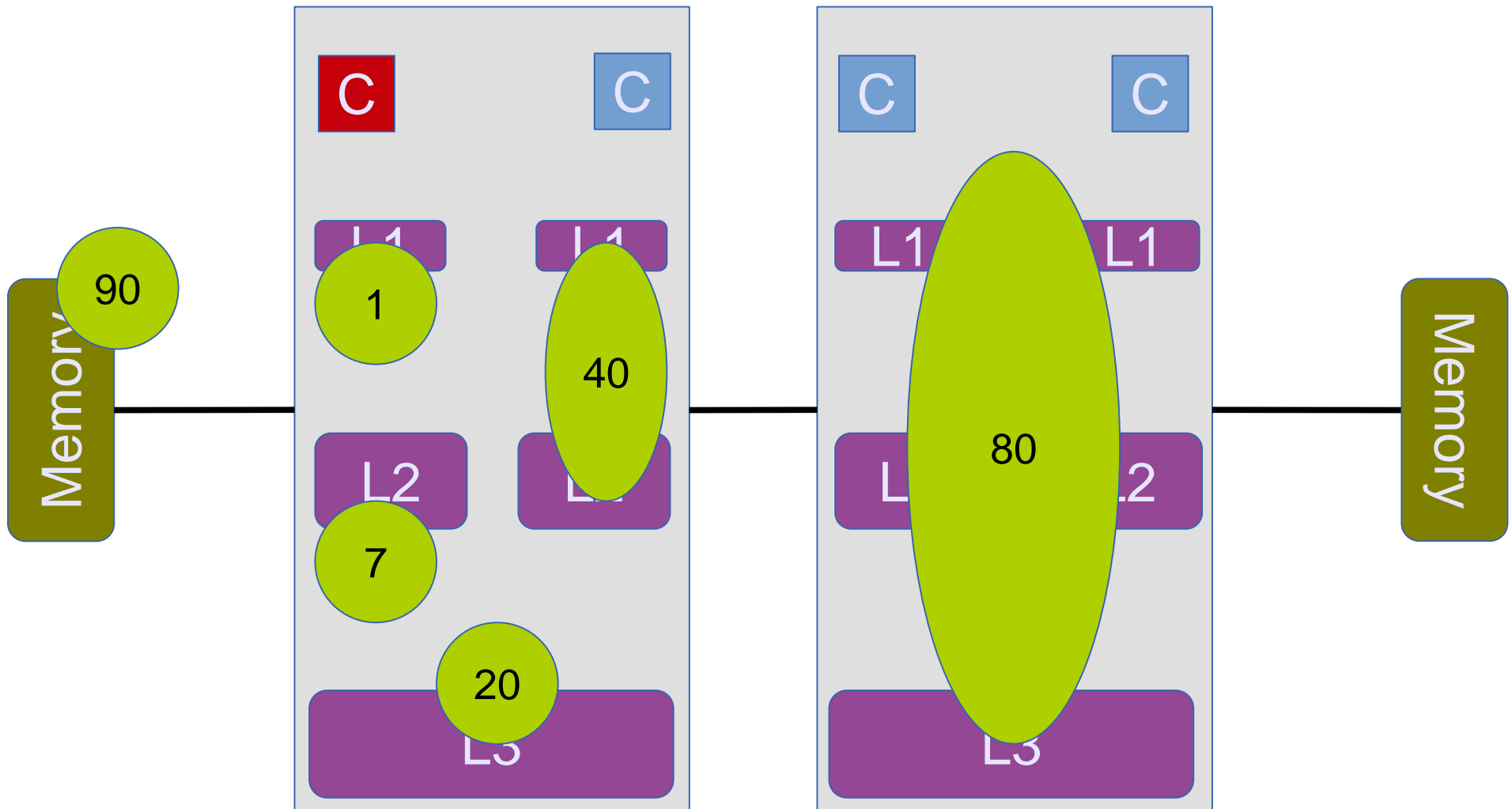
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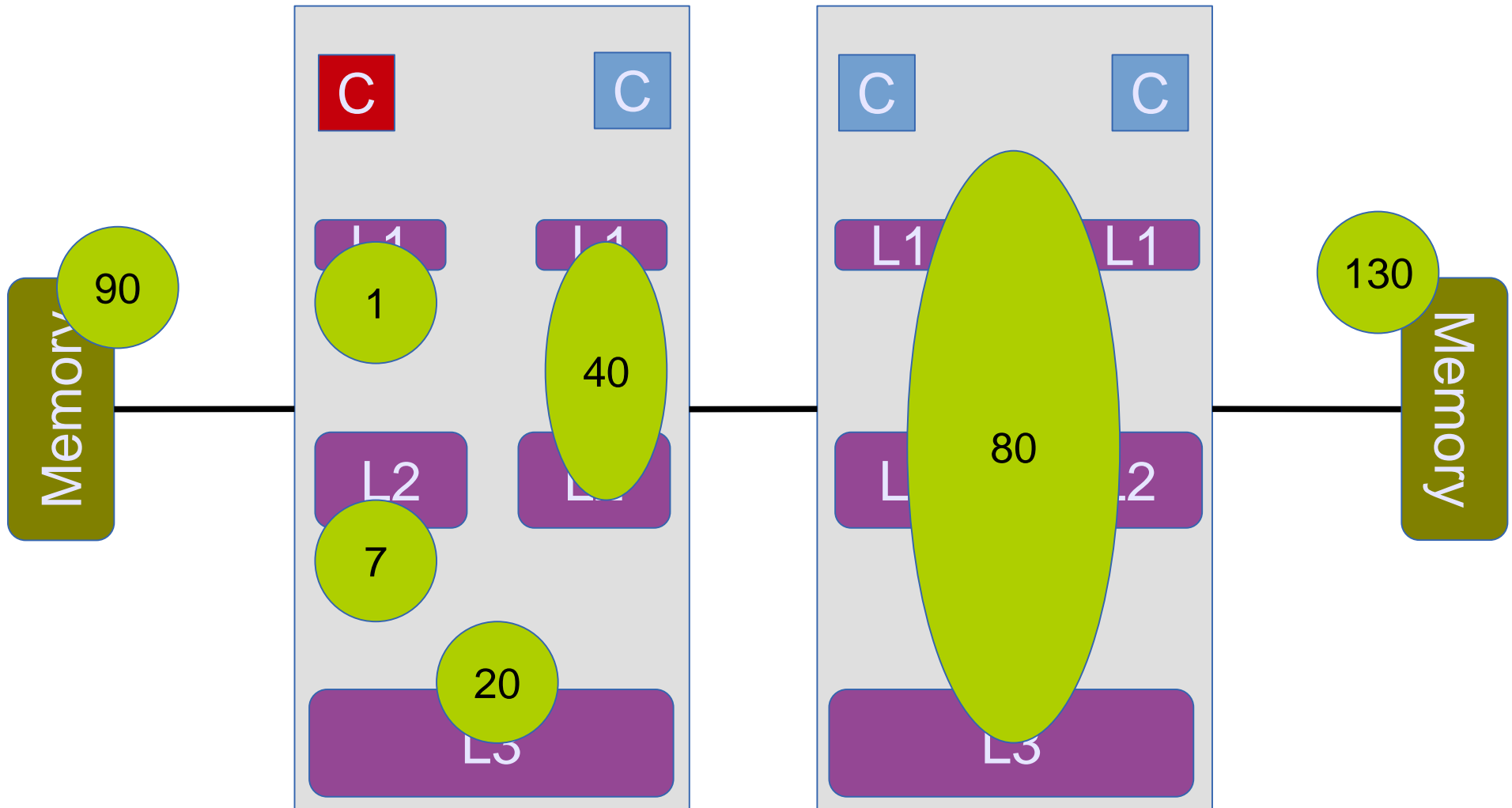
Latency (ns) to access data



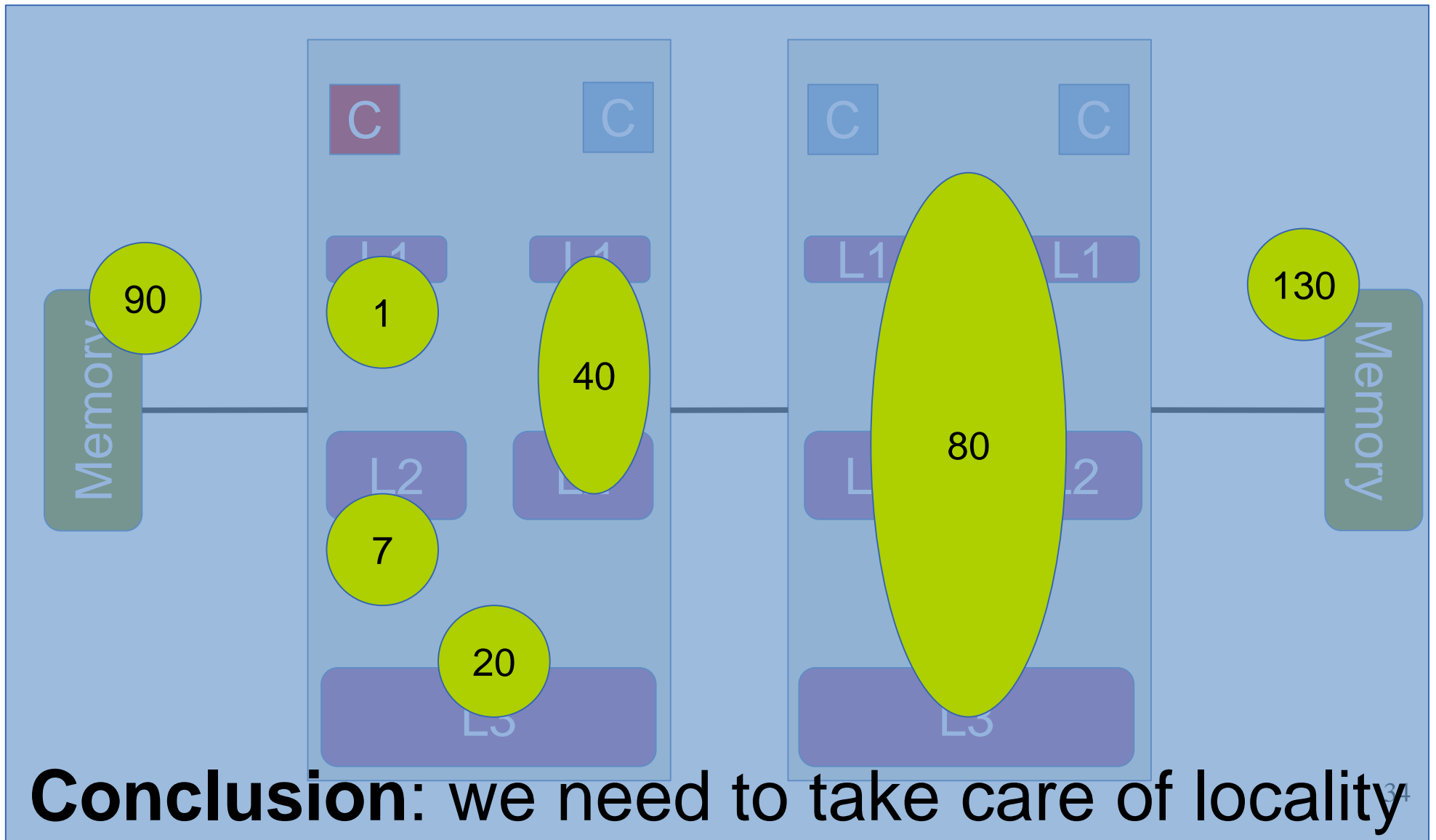
Latency (ns) to access data



Latency (ns) to access data



Latency (ns) to access data



Experiment

The effects of locality

Experiment

The effects of locality

```
vtrigona $ ./test_locality -x0 -y1  
Size:          8 counters = 1 cache lines  
Thread 0 on core : 0  
Thread 1 on core : 2  
Number of threads: 2  
Throughput     : 104.27 Mop/s
```

```
vtrigona $ ./test_locality -x0 -y10  
Size:          8 counters = 1 cache lines  
Thread 0 on core : 0  
Thread 1 on core : 10  
Number of threads: 2  
Throughput     : 43.16 Mop/s
```

Same memory node

Different memory nodes

Outline

- CPU caches
- Cache coherence
- Placement of data
- **Graph processing: Concurrent data structures**

Graph processing

Relational view

People
Table

Name	Likes
Vasilis	Breaking bad
Rachid	Dexter
Vasilis	Dexter

Series
Table

Name	Similar
Breaking bad	Dexter
Dexter	Breaking bad

Graph processing

Relational view

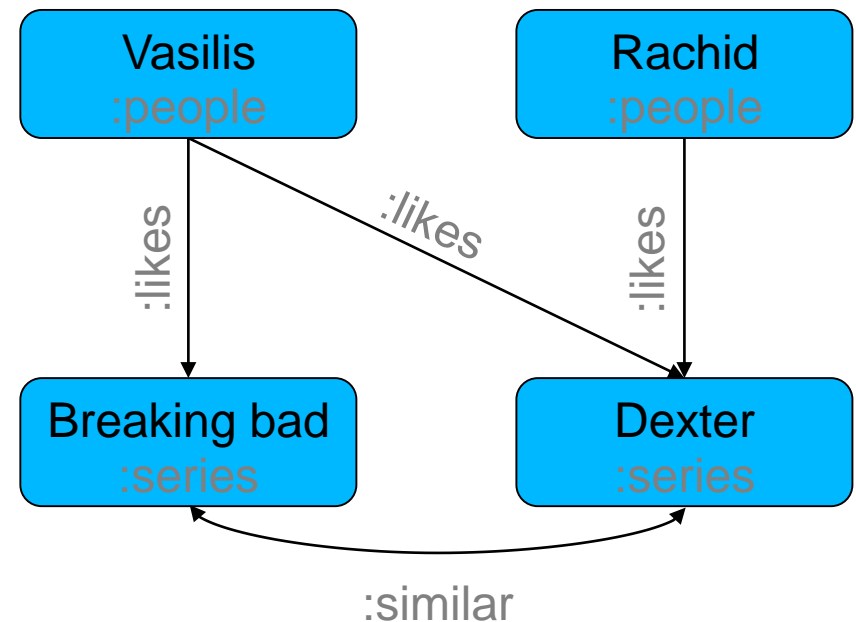
People
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Dexter	Breaking bad

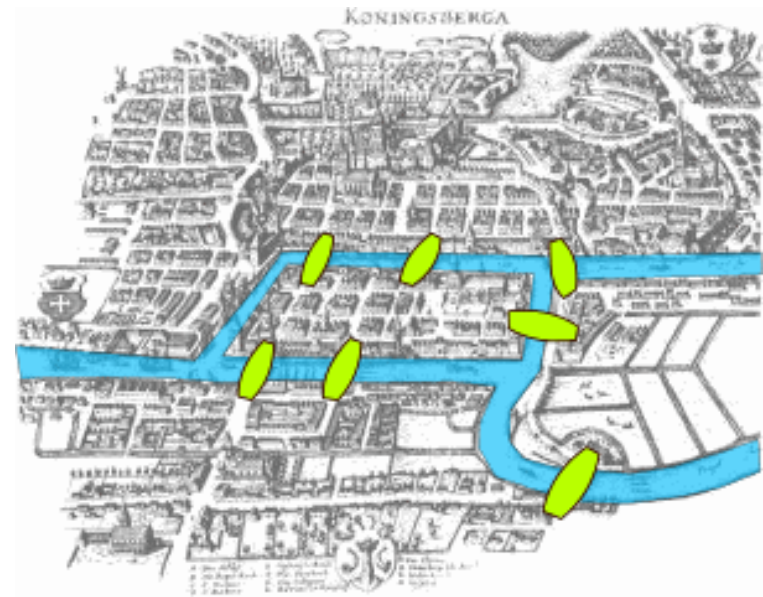
Graph view



Graphs keep the connections among entities materialized

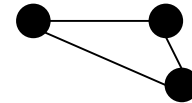
Graph analytics

- Graphs have been studied in Math for centuries
 - Since Euler's "Seven Bridges of Königsberg", 1736
- **Repeatedly traverse your graph and calculate math properties**
- Classic graph problems
 - Graph isomorphism
 - Travelling salesman's problem
 - Max flow, min cut
 - ...
- More recent developments
 - Pagerank
 - Infomap



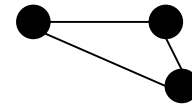
Graph queries

- **Graph pattern matching**
 - Query graphs to find sub-graphs that match a pattern
e.g., triangle counting
- Essentially: SQL for graphs



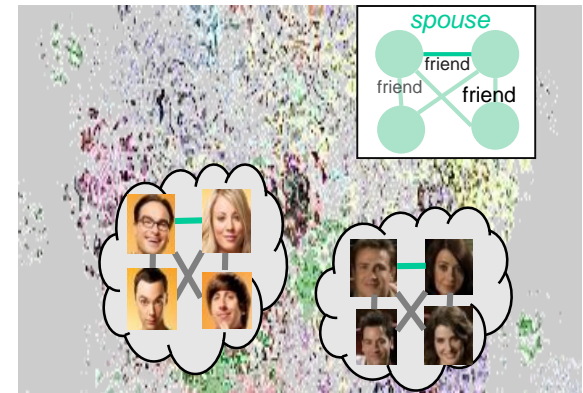
Graph queries

- **Graph pattern matching**
 - Query graphs to find sub-graphs that match a pattern
e.g., triangle counting



- Essentially: SQL for graphs
- Example: Friends of my friends

```
SELECT p1, p3, COUNT(p2)
MATCH (p1)-[:friend]->(p2)->[:friend]->(p3),
      ! (p1)-[:friend]->(p3)
WHERE p1.country = p2.country
GROUP BY p1, p3
ORDER BY COUNT(p2) DESC
```



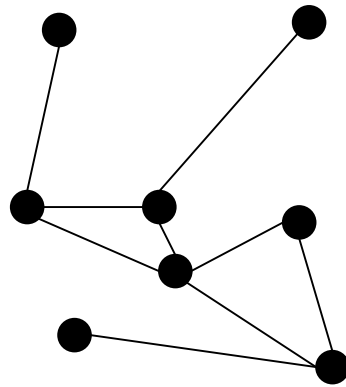
Graph processing frequently involves both analytics and queries

Dissecting a graph processing system

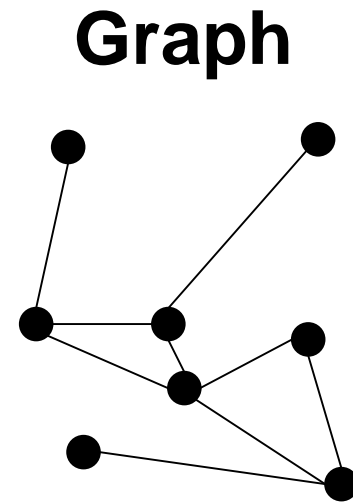
with a focus on (concurrent) data structures

Architecture of a graph processing system

Graph



Architecture of a graph processing system



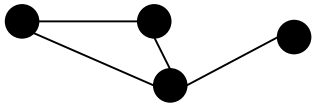
Tons of other data and metadata to store

Graph

tmp graph structure

↓
“Vasilis”, “Breaking bad”, :likes
“Rachid”, “Dexter”, :likes
“Vasilis”, “Dexter”, :likes
“Dexter”, “Breaking bad”, :similar
“Breaking bad”, “Dexter”, :similar

graph structure



user-ids - internal ids

Vasilis → 0	0 → Vasilis
Rachid → 1	1 → Rachid
Breaking bad → 2	2 → Breaking bad
Dexter → 3	3 → Dexter

labels

:likes, :people, :similar, ...

properties

“Vasilis”, {people, male}, 33, Zurich
“Rachid”, {people, male}, ??, Lausanne

lifetime management

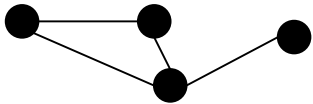
number_of_references: X

Graph

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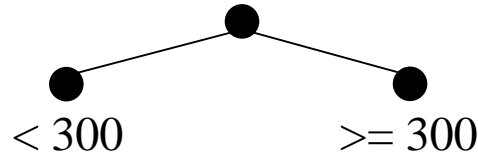
“Vasilis”, {people, male}, 33, Zurich
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lifetime management

number_of_references: X

Runtime

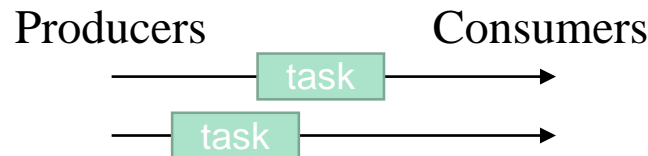
indices / metadata



buffer management

1MB 1MB 1MB 1MB

task / job scheduling



labels

:likes, :people, :similar, :male ...

↓ ↓ ↓ ↓

1 2 3 4

{people, male} → {2,4}

renaming (ids)

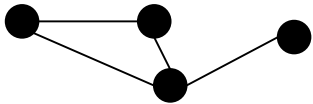
used used used

Graph

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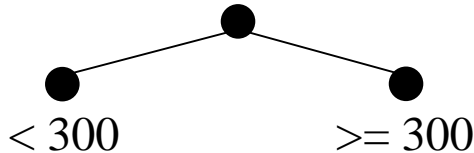
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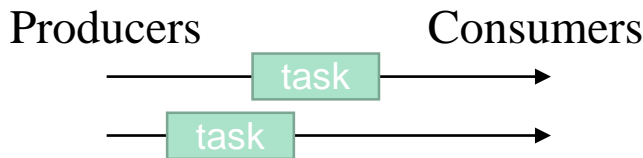
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↓ ↓ ↓ ↓
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renaming (ids)



Operations

group by / join

Vasilis, Breaking bad	→	Vasilis, 2
Rachid, Dexter		Rachid, 1
Vasilis, Dexter		

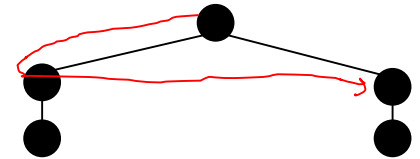
distinct

Vasilis	→	Vasilis
Rachid		Rachid
Vasilis		

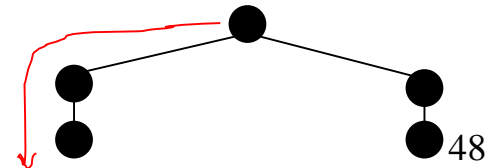
limit (top k)

11 12 0 9 8 13	→	32
8 9 11 23 32 9		23
1 2 3 5 7 3 2 0		13

BFS



DFS

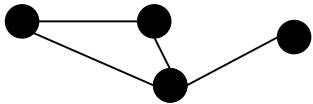


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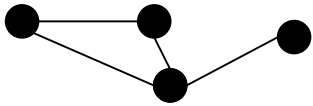
- tmp graph structure
 - append only
 - dynamic schema

Graph

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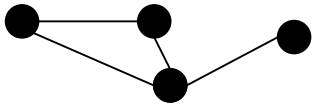
- tmp graph structure
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 - dynamic schema→ **segmented table**

Graph

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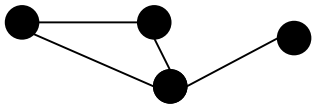
- tmp graph structure
 - append only
 - dynamic schema
 - **segmented table**
- Classic graph structures

Graph

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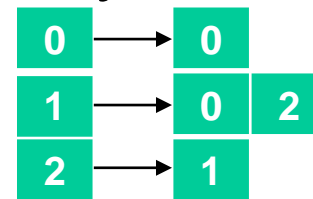
- tmp graph structure
 - append only
 - dynamic schema→ **segmented table**

- Classic graph structures

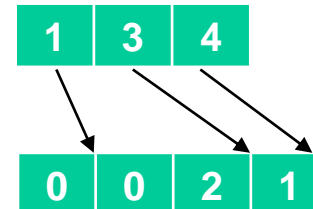
1. connectivity matrix

	0	1	2
0	x		
1	x		x
2		x	

2. adjacency list



3. compressed source row (CSR)



Graph

tmp graph structure

↓
segmented buffer
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“Rachid”, “Dexter”, :likes
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properties

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lifetime management

number_of_references: X

- Mapping user ids to internal ids
 - create once
 - read-only after

Graph

tmp graph structure

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segmented buffer
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Vasilis → 0	0 → Vasilis
Rachid → 1	1 → Rachid
Breaking bad → 2	2 → Breaking bad
Dexter → 3	3 → Dexter

labels

:likes, :people, :similar, ...

properties

“Vasilis”, {people, male}, 33, Zurich
“Rachid”, {people, male}, ??, Lausanne

lifetime management

number_of_references: X

- Mapping user ids to internal ids
 - create once
 - read-only after→ **hash map, lock-free reads**

Graph

tmp graph structure

↓
segmented buffer
“Vasilis”, “Breaking bad”, :likes
“Rachid”, “Dexter”, :likes
“Dexter”, “Breaking bad”, :similar
“Breaking bad”, “Dexter”, :similar

graph structure



user-ids - internal ids

Vasilis → 0	0 → Vasilis
Rachid → 1	1 → Rachid
Breaking bad → 2	2 → Breaking bad
Dexter → 3	3 → Dexter

labels

:likes, :people, :similar, ...

properties

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“Rachid”, {people, male}, ??, Lausanne

lifetime management

number_of_references: X

- Mapping user ids to internal ids
 - create once
 - read-only after
 - **hash map, lock-free reads**
- Mapping internal ids to user ids
 - create once
 - read-only after
 - fixed key range: [0, N}

Graph

tmp graph structure

↓
segmented buffer
“Vasilis”, “Breaking bad”, :likes
“Rachid”, “Dexter”, :likes
“Dexter”, “Breaking bad”, :similar
“Breaking bad”, “Dexter”, :similar

graph structure



user-ids - internal ids

Vasilis → 0	0 → Vasilis
Rachid → 1	1 → Rachid
Breaking bad → 2	2 → Breaking bad
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labels

:likes, :people, :similar, ...

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lifetime management

number_of_references: X

- Mapping user ids to internal ids
 - create once
 - read-only after
 - **hash map, lock-free reads**
- Mapping internal ids to user ids
 - create once
 - read-only after
 - fixed key range: [0, N]
 - **(sequential) array**

Graph

tmp graph structure

↓
segmented buffer
“Vasilis”, “Breaking bad”, :likes
“Rachid”, “Dexter”, :likes
“Dexter”, “Breaking bad”, :similar
“Breaking bad”, “Dexter”, :similar

graph structure



user-ids - internal ids

hash map / array
Vasilis → 0 0 → Vasilis
Rachid → 1 1 → Rachid
Breaking bad → 2 2 → Breaking bad
Dexter → 3 3 → Dexter

labels

:likes, :people, :similar, ...

properties

“Vasilis”, {people, male}, 33, Zurich
“Rachid”, {people, male}, ??, Lausanne

lifetime management

number_of_references: X

- Storing labels
 - usually a small enumeration e.g., person, female, male
 - storing strings is expensive
“person” → ~ 7 bytes
 - comparing strings is expensive

Graph

tmp graph structure

↓
segmented buffer
“Vasilis”, “Breaking bad”, :likes
“Rachid”, “Dexter”, :likes
“Dexter”, “Breaking bad”, :similar
“Breaking bad”, “Dexter”, :similar

graph structure



user-ids - internal ids

hash map/array
Vasilis → 0 0 → Vasilis
Rachid → 1 1 → Rachid
Breaking bad → 2 2 → Breaking bad
Dexter → 3 3 → Dexter

labels

:likes, :people, :similar, ...

properties

“Vasilis”, {people, male}, 33, Zurich
“Rachid”, {people, male}, ??, Lausanne

lifetime management

number_of_references: X

- Storing labels
 - usually a small enumeration e.g., person, female, male
 - storing strings is expensive
“person” → ~ 7 bytes
 - comparing strings is expensive
→ **dictionary encoding**, e.g.,
 - person → 0
 - female → 1
 - male → 2
- Ofc, **hash map** to
 - store those
 - translate during runtime

Graph

tmp graph structure

↓
segmented buffer
“Vasilis”, “Breaking bad”, :likes
“Rachid”, “Dexter”, :likes
“Dexter”, “Breaking bad”, :similar
“Breaking bad”, “Dexter”, :similar

graph structure



user-ids - internal ids

Vasilis → 0 0 → Vasilis
Rachid → 1 1 → Rachid
Breaking bad → 2 2 → Breaking bad
Dexter → 3 3 → Dexter

labels

:likes, :similar, ...

properties

“Vasilis”, {people, male}, 33, Zurich
“Rachid”, {people, male}, ??, Lausanne

lifetime management

number_of_references: X

• Property

- one type per property, e.g., int
- 1:1 mapping with vertices/edges

Graph

tmp graph structure

↓
segmented buffer
“Vasilis”, “Breaking bad”, :likes
“Rachid”, “Dexter”, :likes
“Dexter”, “Breaking bad”, :similar
“Breaking bad”, “Dexter”, :similar

graph structure



user-ids - internal ids

Vasilis → 0 0 → Vasilis
Rachid → 1 1 → Rachid
Breaking bad → 2 2 → Breaking bad
Dexter → 3 3 → Dexter

labels

:likes, :similar, ...

properties

“Vasilis”, {people, male}, 33, Zurich
“Rachid”, {people, male}, ??, Lausanne

lifetime management

number_of_references: X

• Property

- one type per property, e.g., int
 - 1:1 mapping with vertices/edges
- **(sequential) arrays**

Graph

tmp graph structure

↓
segmented buffer
“Vasilis”, “Breaking bad”, :likes
“Rachid”, “Dexter”, :likes
“Dexter”, “Breaking bad”, :similar
“Breaking bad”, “Dexter”, :similar

graph structure



user-ids - internal ids

Vasilis → 0 0 → Vasilis
Rachid → 1 1 → Rachid
Breaking bad → 2 2 → Breaking bad
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labels

:likes, :similar, ...

properties

“Vasilis”, {people, male}, 33, Zurich
“Rachid”, {people, male}, ??, Lausanne

lifetime management

number_of_references: X

- Property
 - one type per property, e.g., int
 - 1:1 mapping with vertices/edges
 - **(sequential) arrays**
- Lifetime management (and other counters)
 - cache coherence: atomic counters can be expensive

Graph

tmp graph structure

↓
segmented buffer
“Vasilis”, “Breaking bad”, :likes
“Rachid”, “Dexter”, :likes
“Dexter”, “Breaking bad”, :similar
“Breaking bad”, “Dexter”, :similar

graph structure



user-ids - internal ids

hash map/array
Vasilis → 0 0 → Vasilis
Rachid → 1 1 → Rachid
Breaking bad → 2 2 → Breaking bad
Dexter → 3 3 → Dexter

labels

dictionary
:likes {people, male}, ...

properties

“Vasilis”, {people, male}, 33, Zurich
“Rachid”, {people, male}, ??, Lausanne

lifetime management

number_of_references: X

- Property
 - one type per property, e.g., int
 - 1:1 mapping with vertices/edges→ **(sequential) arrays**

- Lifetime management (and other counters)

- cache coherence: atomic counters can be expensive

- Two potential solutions

1. **approximate counters**
2. **stripped counters**

Thread local: counter[0] counter[1] counter[2]

```
increment(int by) { counter[my_thread_id] += by; }
```

```
int value() {
```

```
  int sum = 0;
```

```
  for (int i = 0; i < num_threads; i++) { sum += counter[i]; }
```

```
  return sum;
```

```
}
```

Graph

tmp graph structure

“Vasilis”, “Breaking bad”, :likes

“Rachid”, “Dexter”, :likes

“Dexter”, “Breaking bad”, :similar

“Breaking bad”, “Dexter”, :similar

segmented buffer

graph structure



user-ids - internal ids

Vasilis → 0 0 → Vasilis
Rachid → 1 1 → Rachid
Breaking bad → 2 2 → Breaking bad
Dexter → 3 3 → Dexter

hash map / array

labels

dictionary (= map)

properties

“Vasilis”, {people, male}, 33, Zurich
“Rachid”, {people, male}, ??, Lausanne

array

lifetime management

number_of_references: X

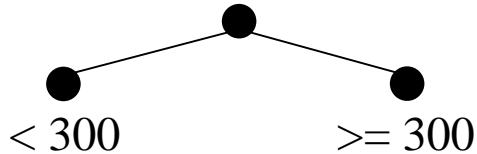
stripped counter

Score

Structure	# Usages
array / buffer	5
map	2

Runtime

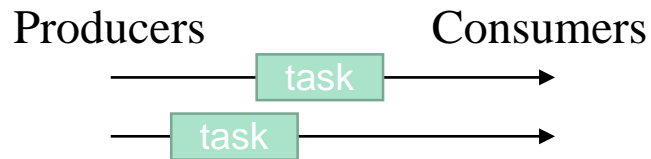
indices / metadata



buffer management

1MB 1MB 1MB 1MB

task / job scheduling



labels

:likes, :people, :similar, :male ...

↓ ↓ ↓ ↓
1 2 3 4

{people, male} → {2,4}

renaming (ids)

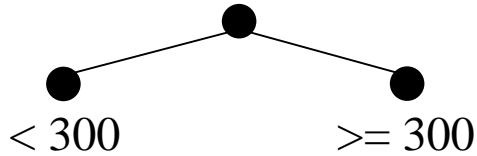
used used used

• Indices

- Used for speeding up “queries”
 - Which vertices have label :person?
 - Which edges have value > 1000?

Runtime

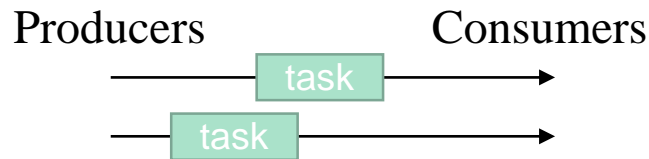
indices / metadata



buffer management

1MB 1MB 1MB 1MB

task / job scheduling



labels

:likes, :people, :similar, :male ...

↓ ↓ ↓ ↓
1 2 3 4

{people, male} → {2,4}

renaming (ids)

used used used

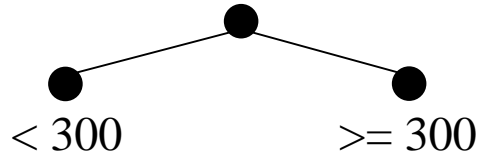
• Indices

- Used for speeding up “queries”
 - Which vertices have label :person?
 - Which edges have value > 1000?

→ **maps, trees**

Runtime

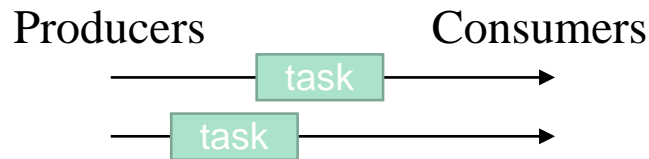
indices / metadata



buffer management

1MB 1MB 1MB 1MB

task / job scheduling



labels

:likes, :people, :similar, :male ...

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renaming (ids)

used used used

• Indices

- Used for speeding up “queries”
 - Which vertices have label :person?
 - Which edges have value > 1000?

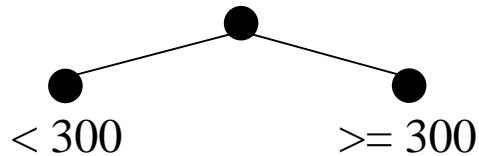
→ **maps, trees**

• Buffer management

- In “real” systems, resource management is very important
- buffer pools
 - no order
 - insertions and deletions
 - no keys

Runtime

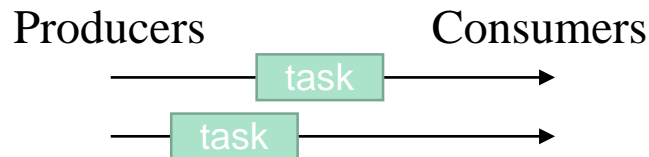
indices / metadata



buffer management

1MB 1MB 1MB 1MB

task / job scheduling



labels

:likes, :people, :similar, :male ...

↓ ↓ ↓ ↓
1 2 3 4

{people, male} → {2,4}

renaming (ids)

used used used

• Indices

- Used for speeding up “queries”
 - Which vertices have label :person?
 - Which edges have value > 1000?

→ **maps, trees**

• Buffer management

- In “real” systems, resource management is very important
- buffer pools
 - no order
 - insertions and deletions
 - no keys

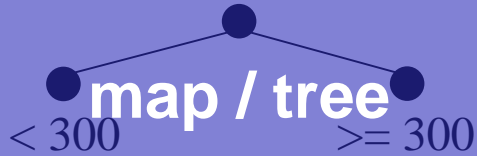
→ Fixed num object pool: **array**

→ Otherwise: **list**

→ Variable-sized elements: **heap**

Runtime

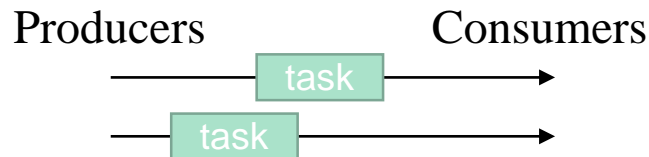
indices / metadata



buffer management



task / job scheduling



labels

:likes, :people, :similar, :male ...



{people, male} → {2,4}

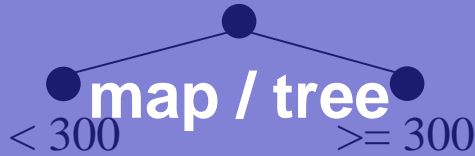
renaming (ids)

used used used

- Task and job scheduling
 - producers create and share tasks
 - consumers get and handle tasks
 - insertions and deletions
 - usually FIFO requirements

Runtime

indices / metadata

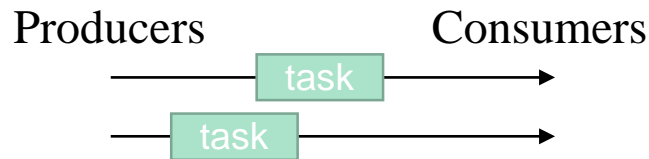


buffer management

array

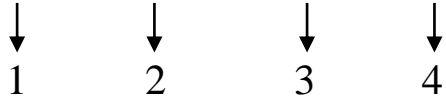


task / job scheduling



labels

:likes, :people, :similar, :male ...



{people, male} → {2,4}

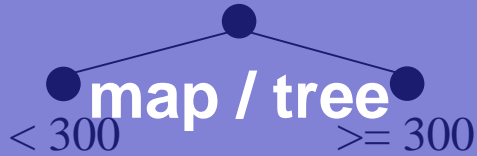
renaming (ids)

used used used

- Task and job scheduling
 - producers create and share tasks
 - consumers get and handle tasks
 - insertions and deletions
 - usually FIFO requirements→ **queues**
- Storing / querying sets of labels
 - set equality expensive
 - usually common groups
e.g., {person, female}, {person, male}

Runtime

indices / metadata

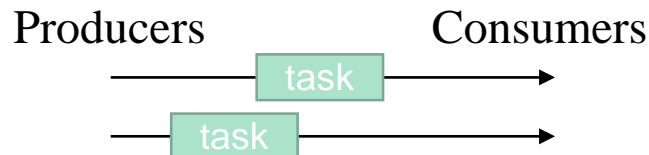


buffer management

array



task / job scheduling



labels

:likes, :people, :similar, :male ...



{people, male} → {2,4}

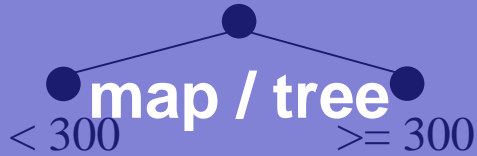
renaming (ids)

used used used

- Task and job scheduling
 - producers create and share tasks
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 - insertions and deletions
 - usually FIFO requirements→ **queues**
- Storing / querying sets of labels
 - set equality expensive
 - usually common groups
e.g., {person, female}, {person, male}→ 2-level **dictionary** encoding
 - {person, female} → 0
 - {person, male} → 1

Runtime

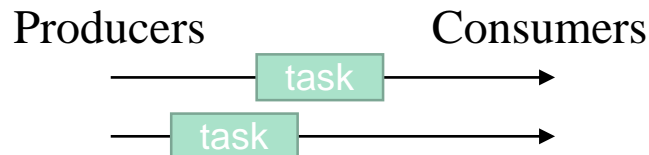
indices / metadata



buffer management



task / job scheduling



labels

:likes, :people, :similar, :male ...



{people, male} → {2,4}

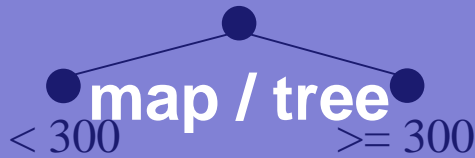
renaming (ids)



- Task and job scheduling
 - producers create and share tasks
 - consumers get and handle tasks
 - insertions and deletions
 - usually FIFO requirements→ **queues**
- Storing / querying sets of labels
 - set equality expensive
 - usually common groups
e.g., {person, female}, {person, male}→ 2-level **dictionary** encoding
 - {person, female} → 0
 - {person, male} → 1
- Giving unique ids (renaming)

Runtime

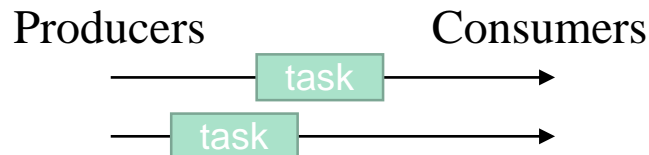
indices / metadata



buffer management

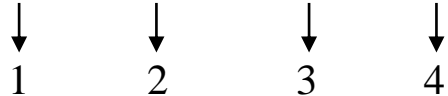


task / job scheduling



labels

:likes, :people, :similar, :male ...



{people, male} → {2,4}

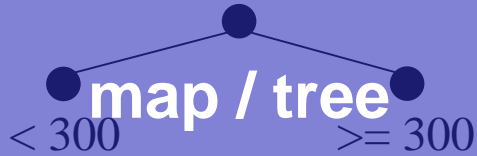
renaming (ids)

used used used

- Task and job scheduling
 - produces create and share tasks
 - consumers get and handle tasks
 - insertions and deletions
 - usually FIFO requirements→ **queues**
- Storing / querying sets of labels
 - set equality expensive
 - usually common groups
e.g., {person, female}, {person, male}→ 2-level **dictionary** encoding
 - {person, female} → 0
 - {person, male} → 1
- Giving unique ids (renaming)
→ **tree, map, set, counter, other?**

Runtime

indices / metadata



buffer management



task / job scheduling



labels

:likes, :people, :similar, :male ...

dictionary (= map)

{people, male} → {2,4}

renaming (ids)

used map / tree / set used

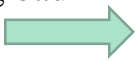
Score

Structure	# Usages
array / buffer	6
map	5
tree / heap	2
set	1
queue	1

Operations

group by / join

Vasilis, Breaking bad
Rachid, Dexter
Vasilis, Dexter



Vasilis, 2
Rachid, 1

distinct


Vasilis
Rachid
Vasilis



Vasilis
Rachid

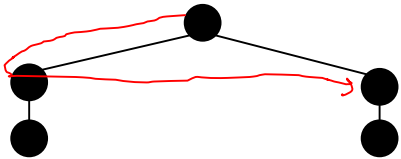
limit (top k)

11 12 0 9 8 13
8 9 11 23 32 9
1 2 3 5 7 3 2 0

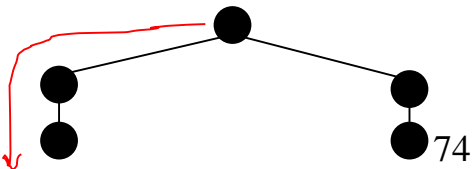


32
23
13

BFS



DFS

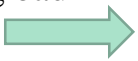


- Group by
 1. Mapping from keys to values
 2. Atomic value aggregations
e.g., COUNT, SUM, MAX
- insertion only

Operations

group by / join

Vasilis, Breaking bad
Rachid, Dexter
Vasilis, Dexter



Vasilis, 2
Rachid, 1

distinct


Vasilis
Rachid
Vasilis



Vasilis
Rachid

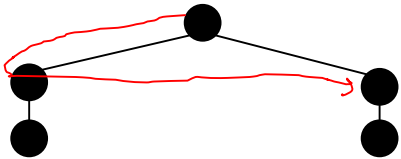
limit (top k)

11 12 0 9 8 13
8 9 11 23 32 9
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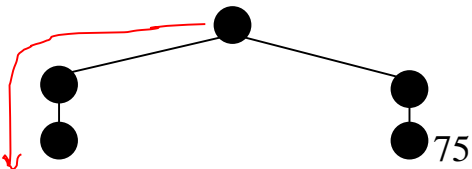


32
23
13

BFS



DFS

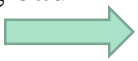


- Group by
 1. Mapping from keys to values
 2. Atomic value aggregations
e.g., COUNT, SUM, MAX
 - insertion only
 - hash map
 - atomic inc / sum / max, etc.

Operations

group by / join

Vasilis, Breaking bad
Rachid, Dexter
Vasilis, Dexter



Vasilis, 2
Rachid, 1

distinct


Vasilis
Rachid
Vasilis



Vasilis
Rachid

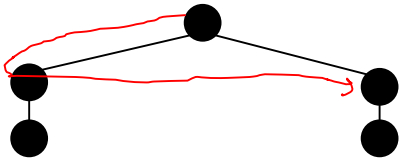
limit (top k)

11 12 0 9 8 13
8 9 11 23 32 9
1 2 3 5 7 3 2 0

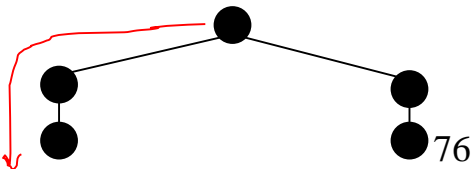


32
23
13

BFS



DFS



- Group by
 1. Mapping from keys to values
 2. Atomic value aggregations
e.g., COUNT, SUM, MAX
 - insertion only


→ hash map

→ atomic inc / sum / max, etc.
- Join
 - create a map of the small table
 - insertion phase, followed by
 - probing phase

Operations

group by / join

Vasilis, Breaking bad
Rachid, Dexter
Vasilis, Dexter



Vasilis, 2
Rachid, 1

distinct


Vasilis
Rachid
Vasilis



Vasilis
Rachid

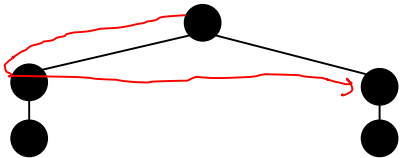
limit (top k)

11 12 0 9 8 13
8 9 11 23 32 9
1 2 3 5 7 3 2 0

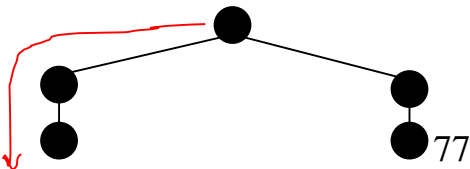


32
23
13

BFS



DFS



- Group by
 1. Mapping from keys to values
 2. Atomic value aggregations
e.g., COUNT, SUM, MAX
 - insertion only

→ hash map

→ atomic inc / sum / max, etc.
- Join
 - create a map of the small table
 - insertion phase, followed by
 - probing phase

→ hash map, lock-free probing

Operations

group by / join

Vasilis, Breaking bad → Vasilis, 2
Rachid, Dexter → Rachid, 1
Vasilis, Dexter → Vasilis, 1

map / atomics

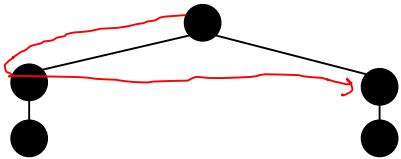
distinct

Vasilis → Vasilis
Rachid → Rachid
Vasilis →

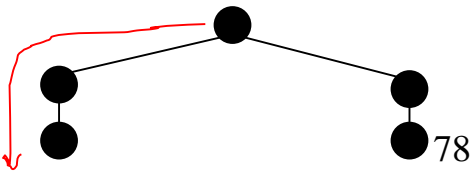
limit (top k)

11 12 0 9 8 13 → 32
8 9 11 23 32 9 → 23
1 2 3 5 7 3 2 0 → 13

BFS



DFS



- Distinct
 - can be solved with sorting, or

Operations

group by / join

Vasilis, Breaking bad → Vasilis, 2
Rachid, Dexter → Rachid, 1
Vasilis, Dexter → Vasilis, 1

map / atomics

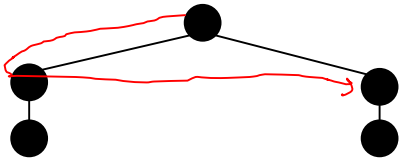
distinct

Vasilis
Rachid
Vasilis → Vasilis
Rachid

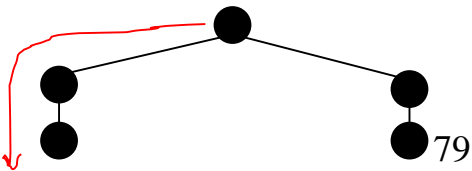
limit (top k)

11 12 0 9 8 13 → 32
8 9 11 23 32 9 → 23
1 2 3 5 7 3 2 0 → 13

BFS



DFS



- Distinct
 - can be solved with sorting, or → **hash set**

Operations

group by / join

Vasilis, Breaking bad
Rachid, Dexter
Vasilis, Dexter

map / atomics

Vasilis, 2
Rachid, 1

distinct

Vasilis
Rachid
Vasilis

→

Vasilis
Rachid

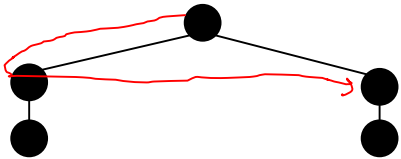
limit (top k)

11 12 0 9 8 13
8 9 11 23 32 9
1 2 3 5 7 3 2 0

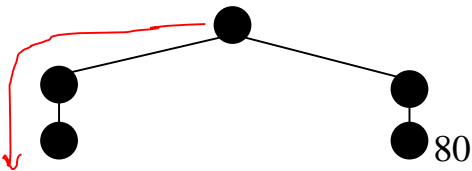
→

32
23
13

BFS



DFS



- Distinct
 - can be solved with sorting, or
→ **hash set**
- Limit (top k)
 - can be solved with sorting, or
 - different specialized structures

Operations

group by / join

Vasilis, Breaking bad
Rachid, Dexter
Vasilis, Dexter

map / atomics

Vasilis, 2
Rachid, 1

distinct

Vasilis
Rachid
Vasilis

→

Vasilis
Rachid

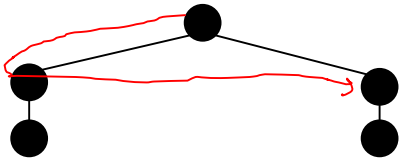
limit (top k)

11 12 0 9 8 13
8 9 11 23 32 9
1 2 3 5 7 3 2 0

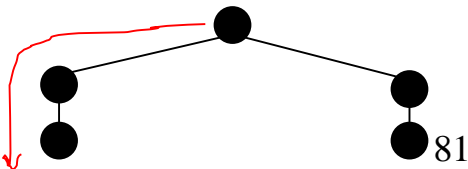
→

32
23
13

BFS



DFS



- Distinct
 - can be solved with sorting, or
→ **hash set**
- Limit (top k)
 - can be solved with sorting, or
 - different specialized structures
→ **tree**
→ **heap**
→ **~ list**
→ **array** (e.g., 2 elements only)
→ **register** (1 element only)

Operations

group by / join

Vasilis, Breaking bad → Vasilis, 2
Rachid, Dexter → Rachid, 1
Vasilis, Dexter → Vasilis, 1

map / atomics

distinct

Vasilis → Vasilis
Rachid → Rachid
Vasilis → Rachid

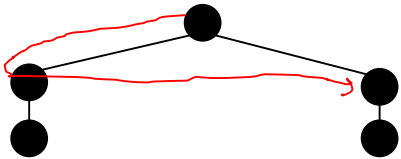
hash set

limit (top k)

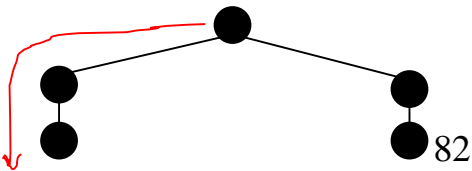
11 12 0 9 8 13 → 32
8 9 11 22 32 9 → 23
1 2 3 5 7 3 2 0 → 13

tree / heap / list

BFS

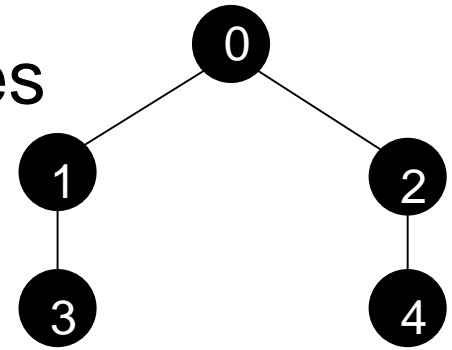


DFS



Breadth-first search (BFS)

- FIFO order
- track visited vertices



Operations

group by / join

Vasilis, Breaking bad → Vasilis, 2
Rachid, Dexter → Rachid, 1
Vasilis, Dexter → Vasilis, 1

map / atomics

distinct

Vasilis → Vasilis
Rachid → Rachid
Vasilis → Rachid

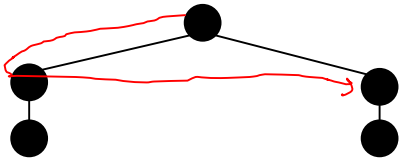
hash set

limit (top k)

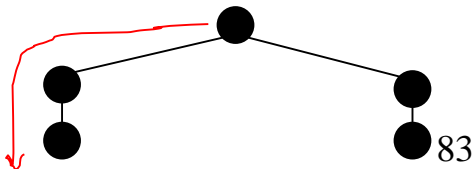
11 12 0 9 8 13 → 32
8 9 11 22 32 9 → 23
1 2 3 5 7 3 2 0 → 13

tree / heap / list

BFS

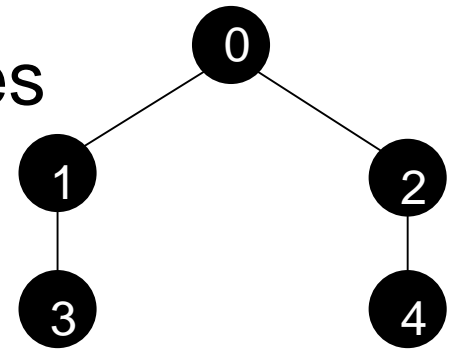


DFS



Breadth-first search (BFS)

- FIFO order
 - track visited vertices
- **queue**
- **set**



Operations

group by / join

Vasilis, Breaking bad → Vasilis, 2
Rachid, Dexter → Rachid, 1
Vasilis, Dexter → Vasilis, 1

map / atomics

distinct

Vasilis → Vasilis
Rachid → Rachid
Vasilis → Rachid

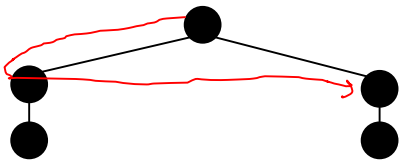
hash set

limit (top k)

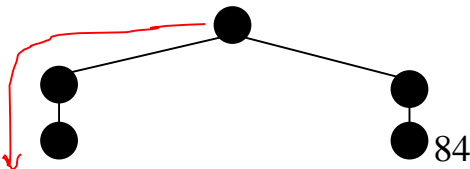
11 12 0 9 8 13 → 32
8 9 11 22 32 9 → 23
1 2 3 5 7 3 2 0 → 13

tree / heap / list

BFS

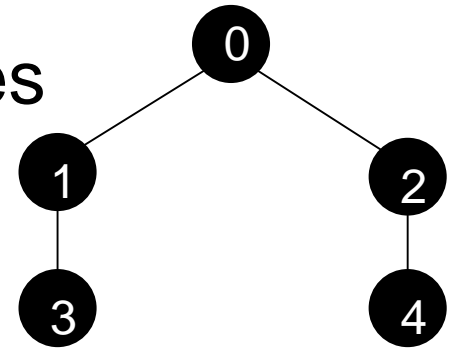


DFS



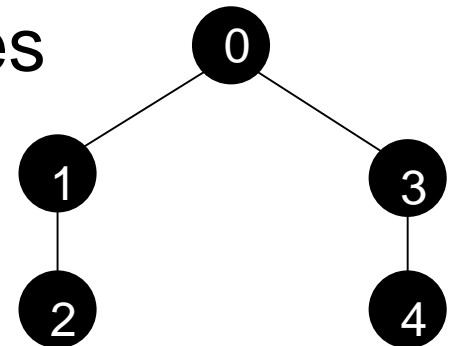
Breadth-first search (BFS)

- FIFO order
 - track visited vertices
- **queue**
→ **set**



Depth-first search (DFS)

- LIFO order
- track visited vertices



Operations

group by / join

Vasilis, Breaking bad → Vasilis, 2
Rachid, Dexter → Rachid, 1
Vasilis, Dexter → Vasilis, 1

map / atomics

distinct

Vasilis
Rachid
Vasilis → Vasilis
Rachid

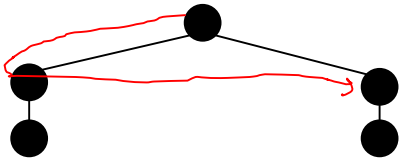
hash set

limit (top k)

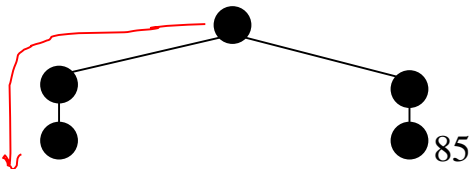
11 12 0 9 8 13 → 32
8 9 11 22 32 9 → 23
1 2 3 5 7 3 20 → 13

tree / heap / list

BFS

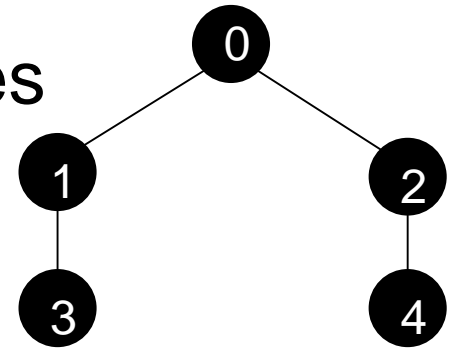


DFS



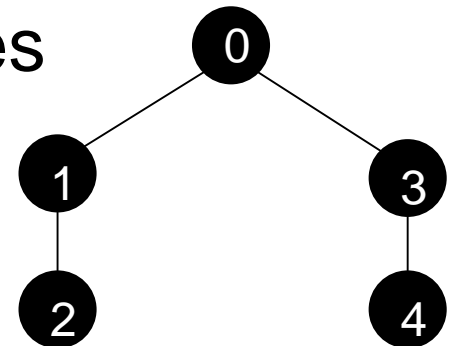
• Breadth-first search (BFS)

- FIFO order
 - track visited vertices
- **queue**
- **set**



• Depth-first search (DFS)

- LIFO order
 - track visited vertices
- **stack**
- **set**



Operations

group by / join

Vasilis, Breaking bad
Rachid, Dexter
Vasilis, Dexter

Vasilis, 2
Rachid, 1

map / atomics

distinct

Vasilis
Rachid
Vasilis

Vasilis
Rachid

hash set

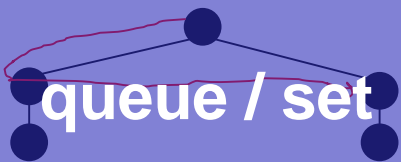
limit (top k)

11 12 0 9 8 13
8 9 11 22 32 9
1 2 3 5 7 3 2 0

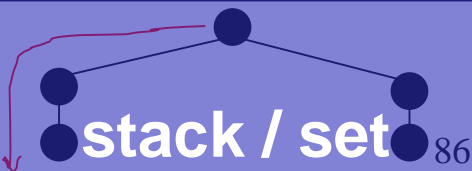
32
23
13

tree / heap / list

BFS



DFS



Score

Structure	# Usages
array / buffer	7
map	6
set	4
tree / heap	3
queue	2
stack	1
list	1

Graph

tmp graph structure

"Vasilis", "Breaking bad", :likes
 "Rachid", "Dexter", :likes
 "Dexter", "Breaking bad", :similar
 "Breaking bad", "Dexter", :similar

segmented buffer

graph structure



user-ids - internal ids

Vasilis → 0 0 → Vasilis
 Rachid → 1 1 → Rachid
 Breaking bad → 2 2 → Breaking bad
 Dexter → 3 3 → Dexter

hash map / array

labels

:likes, :people, :similar, ...

properties

"Vasilis", {people, male}, 33, Zurich
 "Rachid", {people, male}, ??, Lausanne

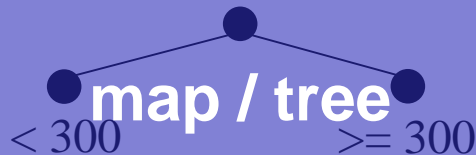
array

lifetime management

number_of_references: X
stripped counter

Runtime

indices / metadata



buffer management



task / job scheduling



labels

:likes, :people, :similar, :male ...



{people, male} → {2,4}

renaming (ids)



Operations

group by / join

Vasilis, Breaking bad
 Rachid, Dexter
 Vasilis, Dexter

map / atomics

distinct

Vasilis
 Rachid
 Vasilis

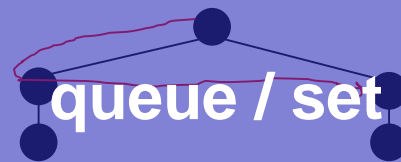
hash set

limit (top k)

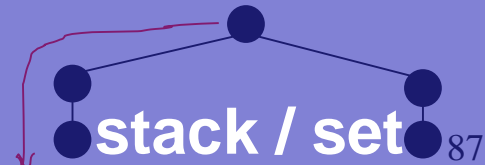
11 12 0 9 8 13
 8 9 11 12 32 9
 1 2 3 5 7 3 2 0

tree / heap / list

BFS



DFS



Conclusions

- Both theory and practice are necessary for
 - Designing, and
 - Implementing fast / scalable data structures
- Hardware plays a huge role on implementations
 - How and which memory access patterns to use
- **(Concurrent) Data structures**
 - **The backbone of every system**
 - **An “open” and challenging area or research**