

# The Power of Registers

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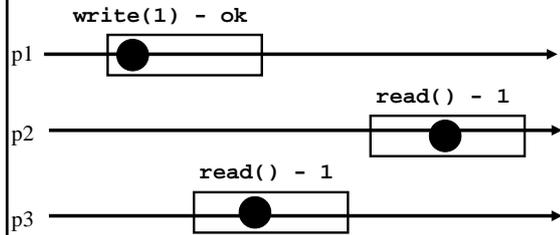


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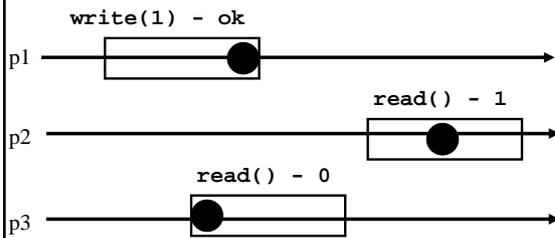


## Atomic execution



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## Atomic execution



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## Registers

- Question 1: what objects can we implement with registers? (this lecture)
- Question 2: what objects we cannot implement? (next lecture)

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## Wait-free implementations of atomic objects

- An **atomic** object is simply defined by its sequential specification; i.e., by how its operations should be implemented when there is no concurrency
- Implementations should be **wait-free**: every process that invokes eventually gets a reply (unless the process crashes)

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## Counter (sequential spec)

- A **counter** has two operations **inc()** and **read()** and maintains an integer  $x$  *init to 0*
- read()**:
  - return( $x$ )
- inc()**:
  - $x := x + 1$ ;
  - return(ok)

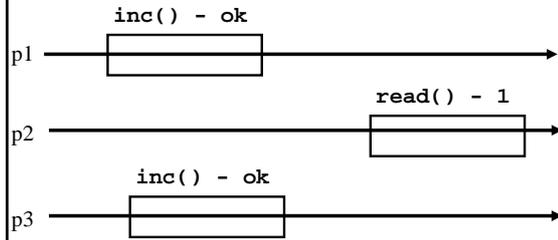
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## Naive implementation

- ☞ The processes share one register Reg
- ☞ *read()*:
  - ☞ return(Reg.read());
- ☞ *inc()*:
  - ☞ temp := Reg.read()+1;
  - ☞ Reg.write(temp);
  - ☞ return(ok)

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## Atomic execution?



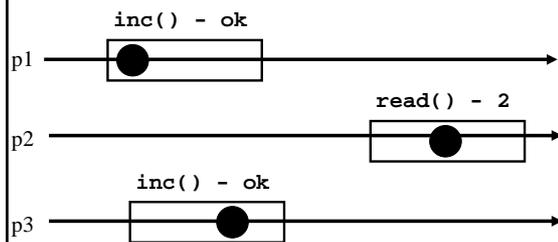
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## Atomic implementation

- ☞ The processes share an array of registers Reg[1,..,n]
- ☞ *inc()*:
  - ☞ temp := Reg[i].read() + 1;
  - ☞ Reg[j].write(temp);
  - ☞ return(ok)

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## Atomic execution?



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## Atomic implementation

- ☞ *read()*:
  - ☞ sum := 0;
  - ☞ for j = 1 to n do
    - ☞ sum := sum + Reg[j].read();
  - ☞ return(sum)

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## Snapshot (sequential spec)

- ☞ A *snapshot* has operations *update()* and *scan()* and maintains an array *x* of size *n*
- ☞ *scan()*:
  - ☞ return(x)
- ☞ *update(i,v)*:
  - ☞ x[i] := v;
  - ☞ return(ok)

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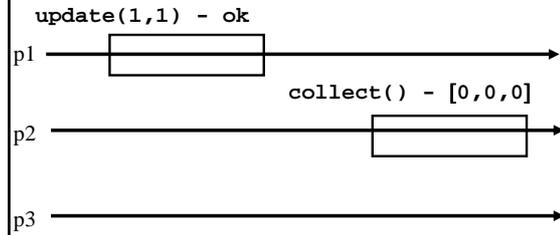
## Very naive implementation

Each process maintains an array of integer variables  $x$  init to  $[0, \dots, 0]$

- **scan()**:
  - return( $x$ )
- **update( $i, v$ )**:
  - $x[i] := v$ ;
  - return(ok)

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## Atomic execution?



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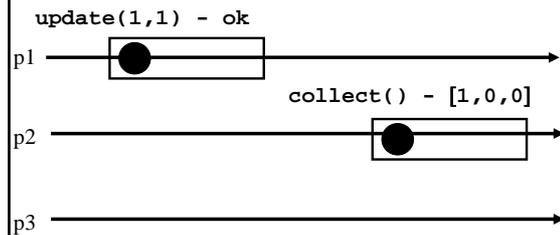
## Less naive implementation

The processes share one array of  $N$  registers  $\text{Reg}[1, \dots, N]$

- **scan()**:
  - for  $j = 1$  to  $N$  do
    - $x[j] := \text{Reg}[j].\text{read}()$ ;
  - return( $x$ )
- **update( $i, v$ )**:
  - $\text{Reg}[i].\text{write}(v)$ ; return(ok)

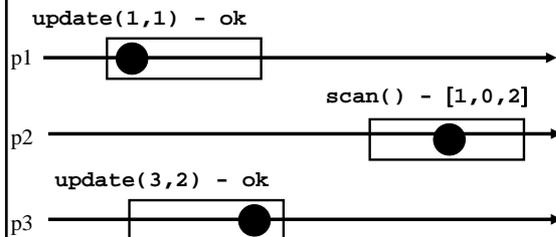
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## Atomic execution?



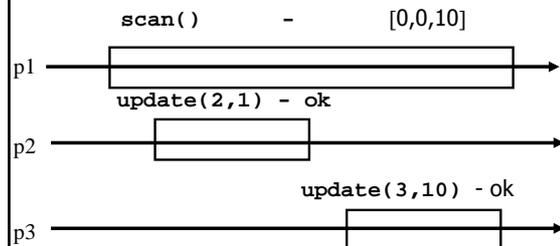
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## Atomic execution?



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## Atomic execution?



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## Non-atomic vs atomic snapshot

- What we implement here is some kind of *regular snapshot*:
- A *scan* returns, for every index of the snapshot, the last written values or the value of any concurrent update
- We call it *collect*

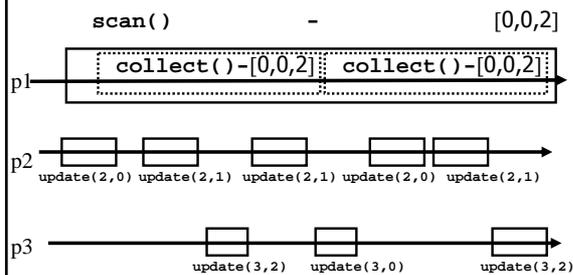
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## Key idea for atomicity

- To *scan*, a process keeps reading the entire snapshot (i.e., it *collect*), until two results at the *same*
- This means that the snapshot did not change, and it is safe to return without violating atomicity

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## Same value vs. Same timestamp



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## Enforcing atomicity

- The processes share one array of  $N$  registers `Reg[1,..,N]`; each contains a value and a timestamp
- We use the following operation for modularity
- collect()*:
  - for  $j = 1$  to  $N$  do
    - `x[j] := Reg[j].read();`
  - return( $x$ )

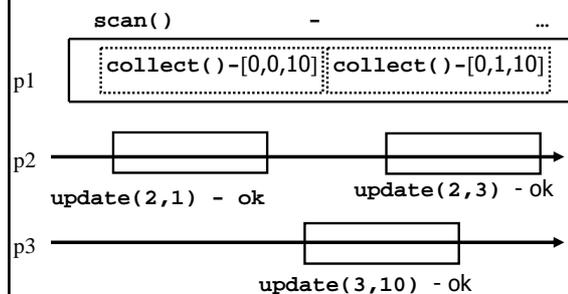
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## Enforcing atomicity (cont'd)

- scan()*:
  - `temp1 := self.collect();`
  - while(true) do
    - `temp2 := self.collect();`
    - `temp1 := temp2;`
    - if (`temp1 = temp2`) then
      - return (`temp1.val`)
- update(i, v)*:
  - `ts := ts + 1;`
  - `Reg[i].write(v, ts);`
  - return(ok)

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## Wait-freedom?



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## Key idea for atomicity & wait-freedom

- ☞ The processes share an array of *registers*  $\text{Reg}[1,\dots,N]$  that contains each:
  - ☞ a value,
  - ☞ a timestamp, and
  - ☞ a copy of the entire array of values

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## Key idea for atomicity & wait-freedom (cont'd)

- ☞ To *scan*, a process keeps collecting and returns a collect if it did not change, or some collect returned by a concurrent *scan*
  - ☞ Timestamps are used to check if the collect changes or if a scan has been taken in the meantime
- To *update*, a process *scans* and writes the value, the new timestamp and the result of the scan

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## Snapshot implementation

Every process keeps a local timestamp  $ts$

- ☞ *update*( $i, v$ ):
  - ☞  $ts := ts + 1$ ;
  - ☞  $\text{Reg}[i].\text{write}(v, ts, \text{self.scan}());$
  - ☞ return(ok)

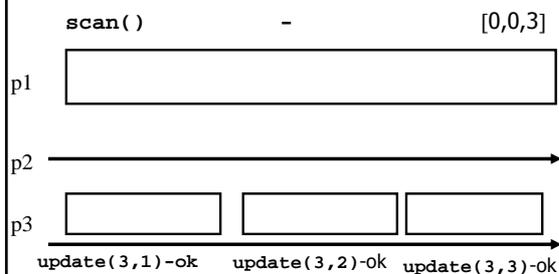
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## Snapshot implementation

- ☞ *scan*():
  - ☞  $t1 := \text{self.collect}(); t2 := t1$
  - ☞ while(true) do
    - ☞  $t3 := \text{self.collect}();$
    - ☞ if ( $t3 = t2$ ) then return ( $t3[j,3]$ );
    - ☞ for  $j = 1$  to  $N$  do
    - ☞ if ( $t3[j,2] \geq t1[j,2] + 2$ ) then
      - ☞ return ( $t3[j,3]$ )
  - ☞  $t2 := t3$

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## Possible execution?



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